

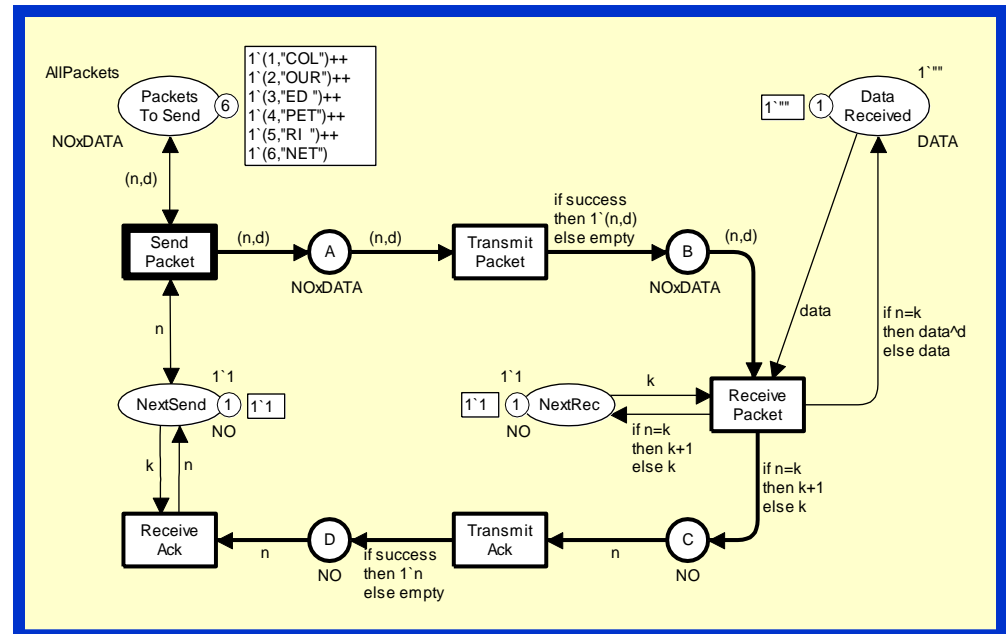
# Coloured Petri Nets

## Modelling and Validation of Concurrent Systems

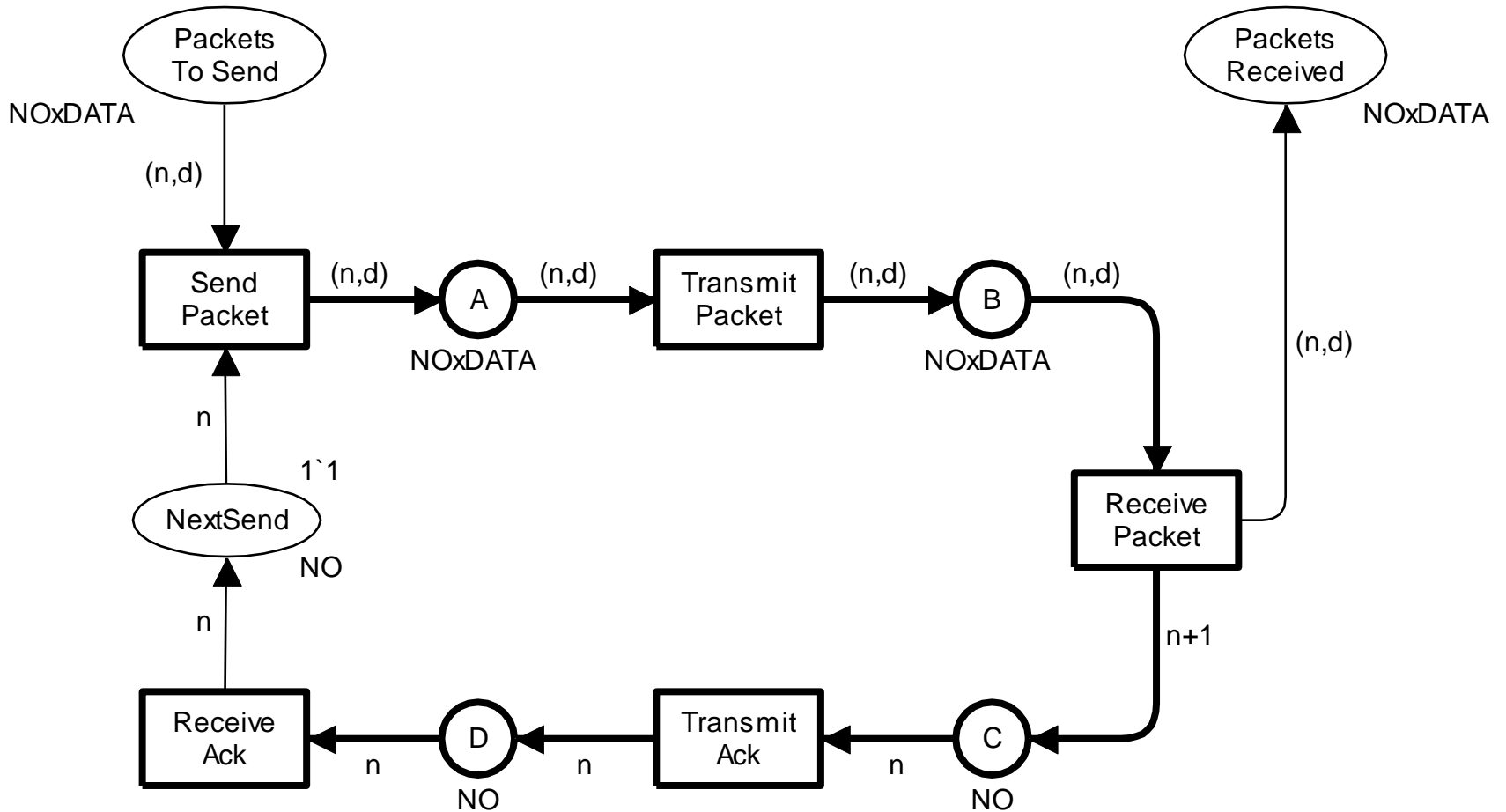
### Chapter 2: Non-hierarchical Coloured Petri Nets

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Lars Michael Kristensen

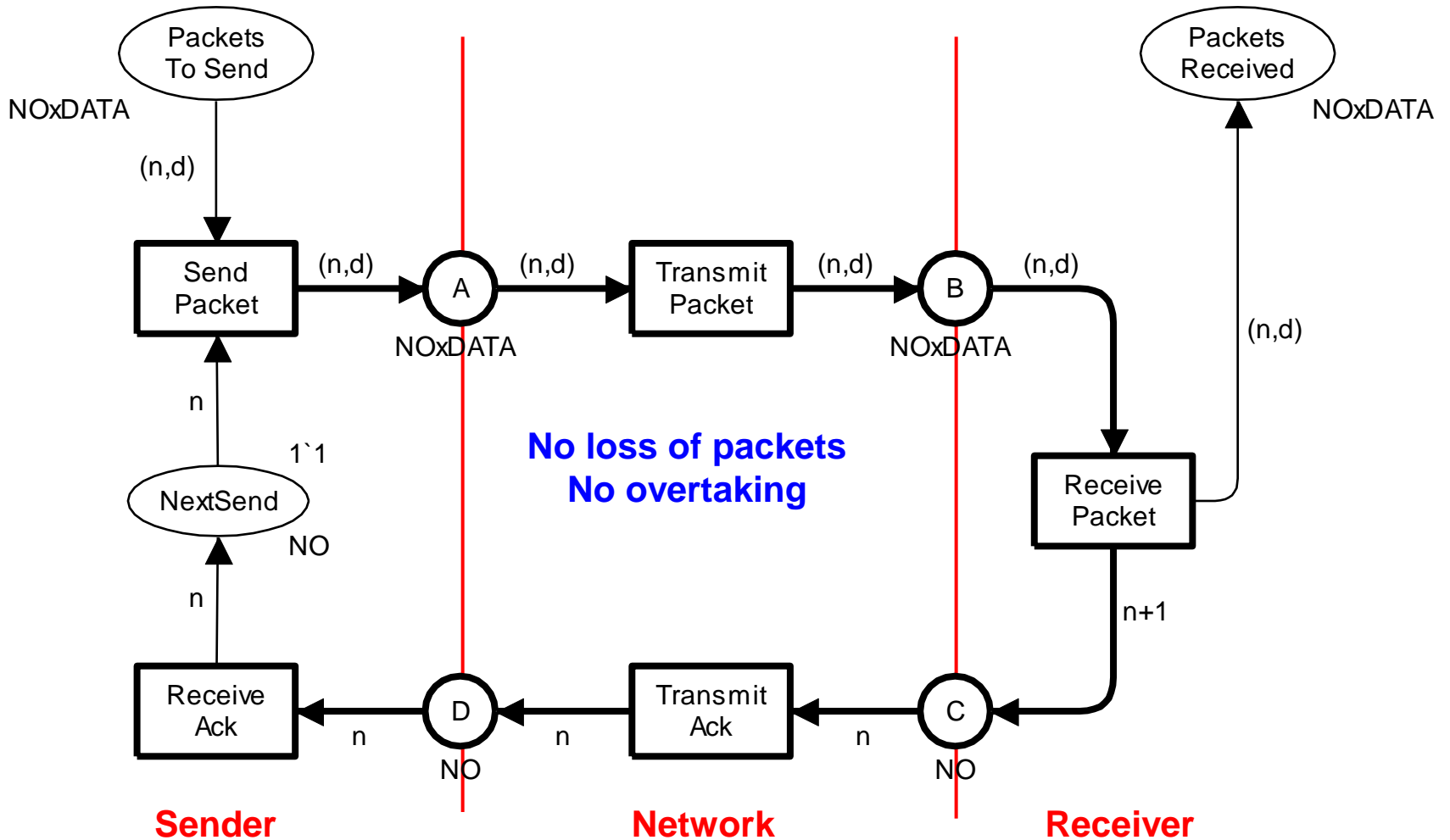
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@cs.au.dk



# Simple protocol

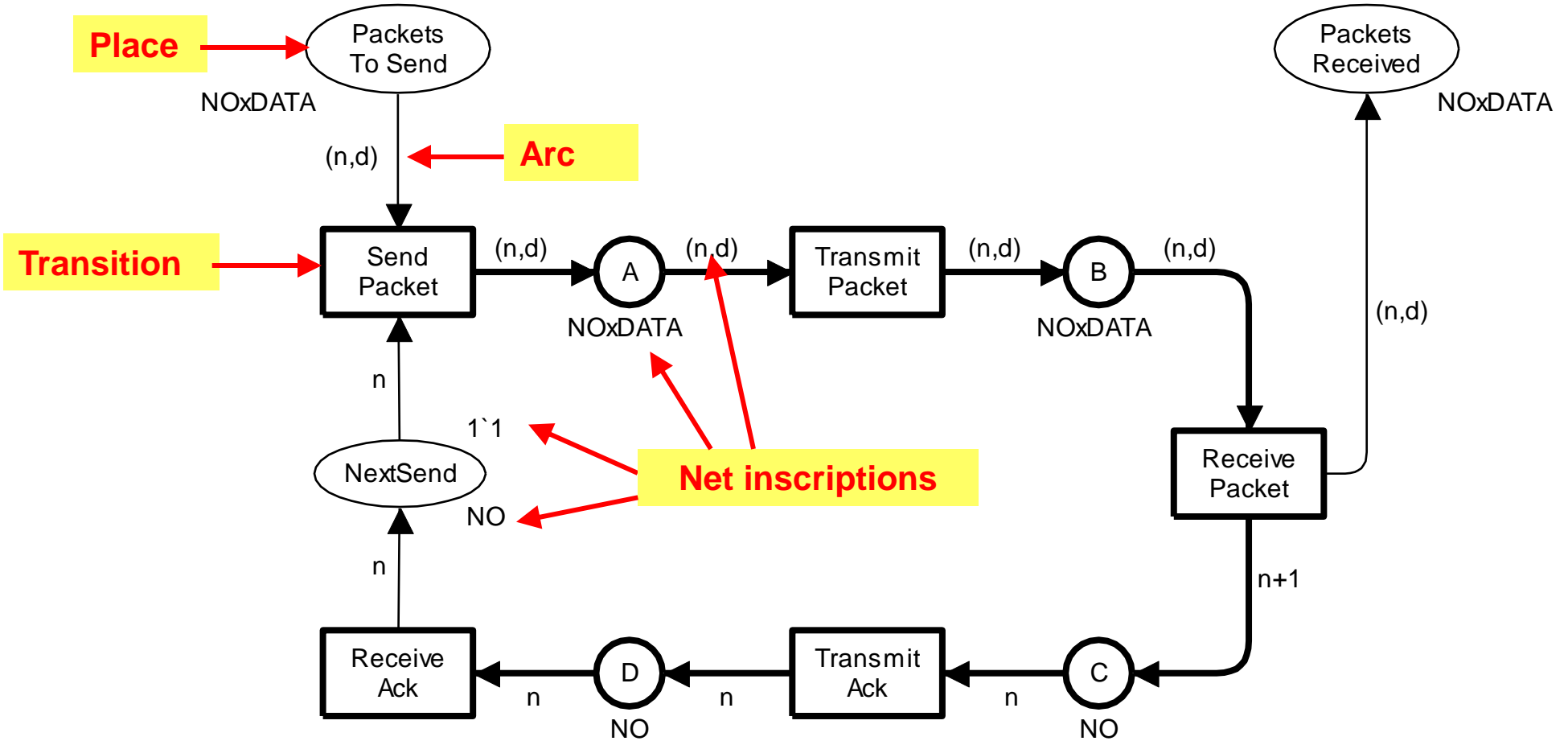


# Informal description



# Coloured Petri Net

$1 \setminus (1, "COL ") ++$   
 $1 \setminus (2, "OUR") ++$   
 $1 \setminus (3, "ED ") ++$   
 $1 \setminus (4, "PET") ++$   
 $1 \setminus (5, "RI ") ++$   
 $1 \setminus (6, "NET")$



# Places represent the state of the system

Initial marking  
(multiset of  
tokens)

Each token in the initial marking  
must have a colour that belongs  
to the colour set

```
1`(1,"COL " )++
1`(2,"OUR" )++
1`(3,"ED " )++
1`(4,"PET" )++
1`(5,"RI " )++
1`(6,"NET" )
```

Name  
(no formal  
meaning;  
large impact  
on readability)

Packets  
To Send

Definition of colour sets:

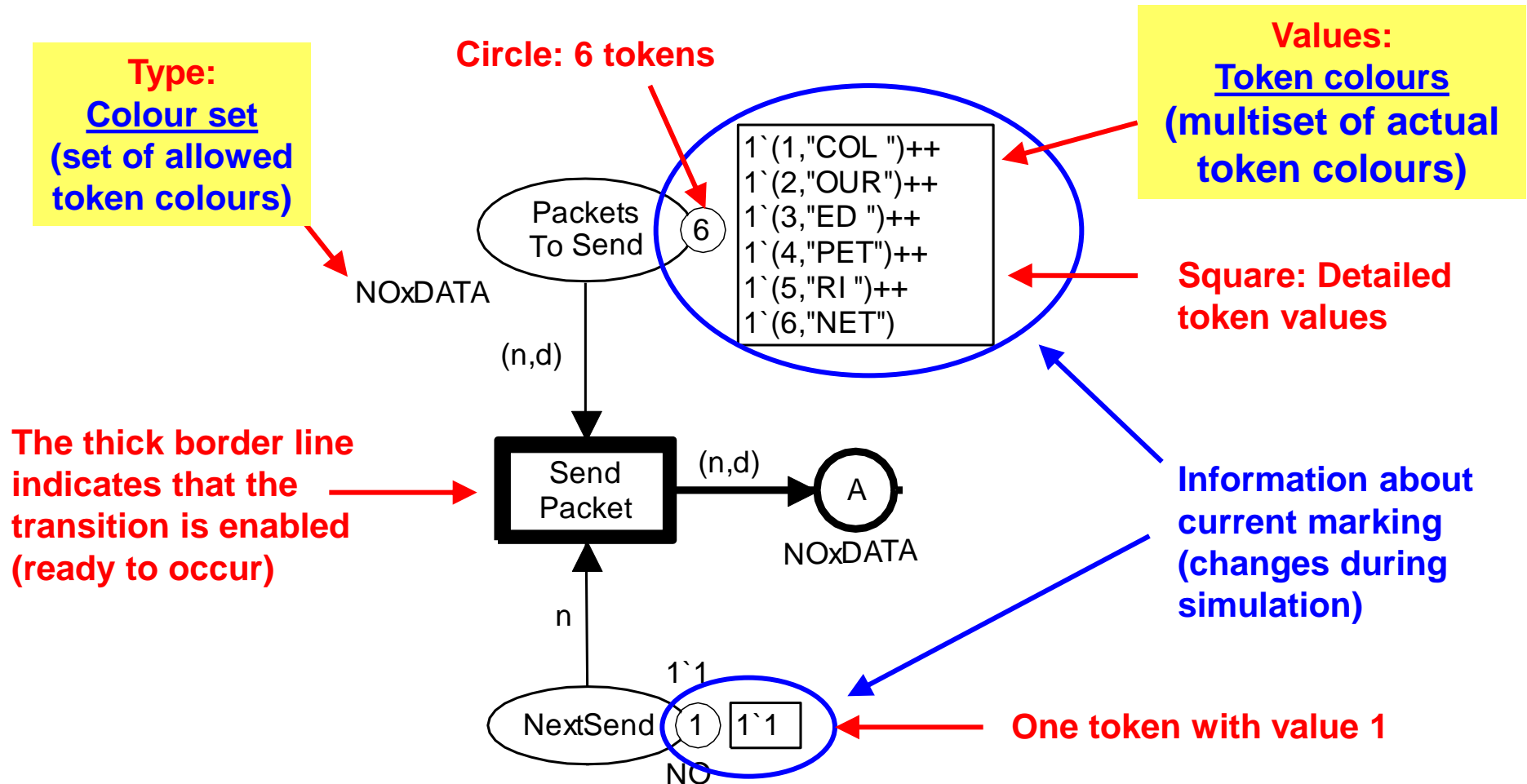
```
colset NO      = int;    (* integers *)
colset DATA   = string; (* text strings *)
colset NOxDATA = product NO * DATA;
```

NOxDATA

Colour set  
(type)

- Each place contains a number of **tokens**.
- Each token carries a **colour** (data value).
- The **colour set** specifies the set of allowed token colours.

# Current marking during simulation



# Transitions and arcs

Arc expression

The type of the arc expression must be equal to the colour set of the attached place (or a multiset over the colour set)

Declaration of variables:

var n : NO; (\* integers \*)  
var d : DATA; (\* strings \*)

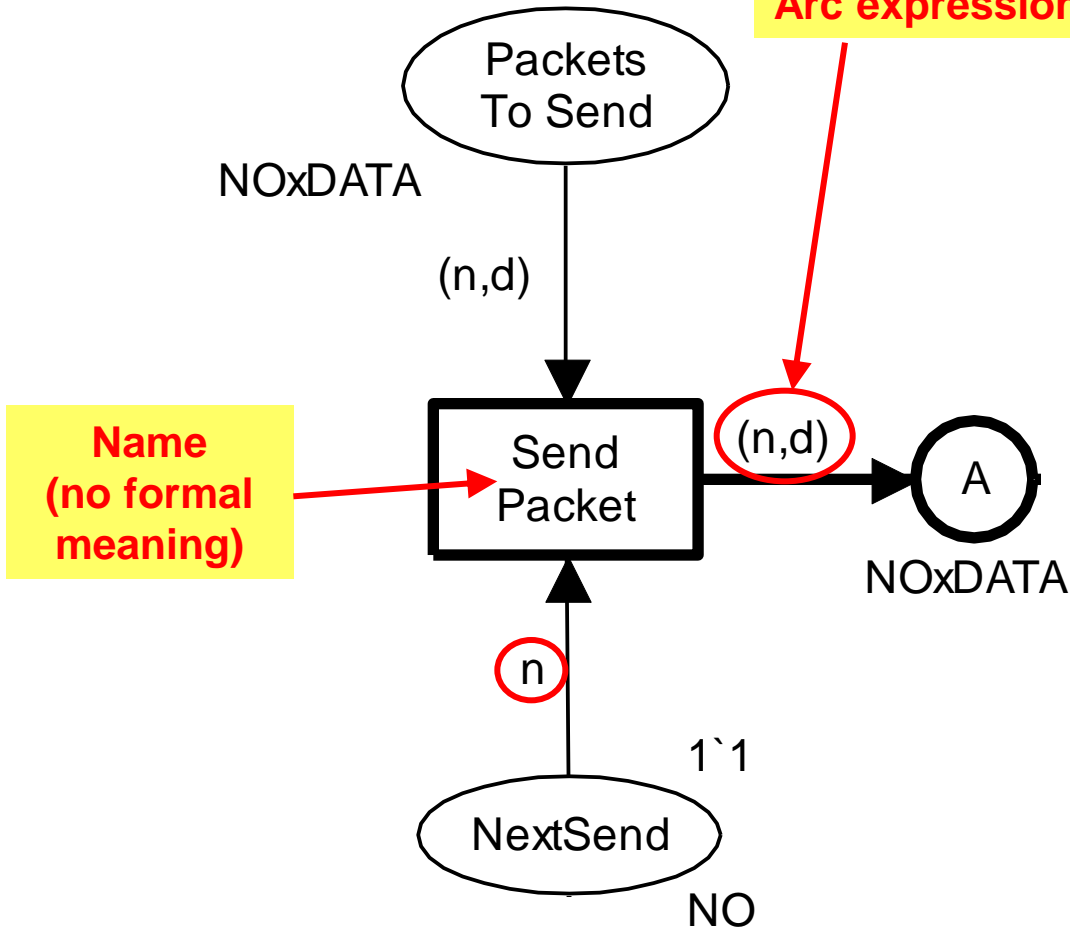
Binding of variables:

<n=3,d="CPN">

Evaluation of expressions:

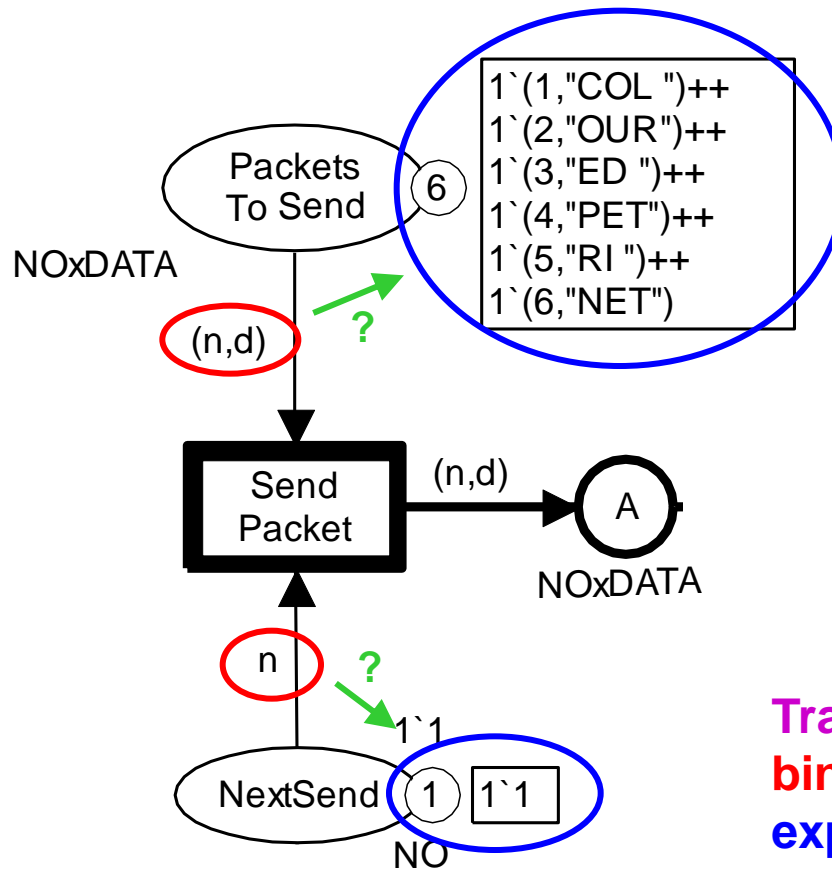
(n,d) → (3,"CPN") : NOxDATA

n → 3 : NO



Name  
(no formal  
meaning)

# Enabling of transitions



Two variables:

var  $n$  : NO; (\* integers \*)  
var  $d$  : DATA; (\* strings \*)

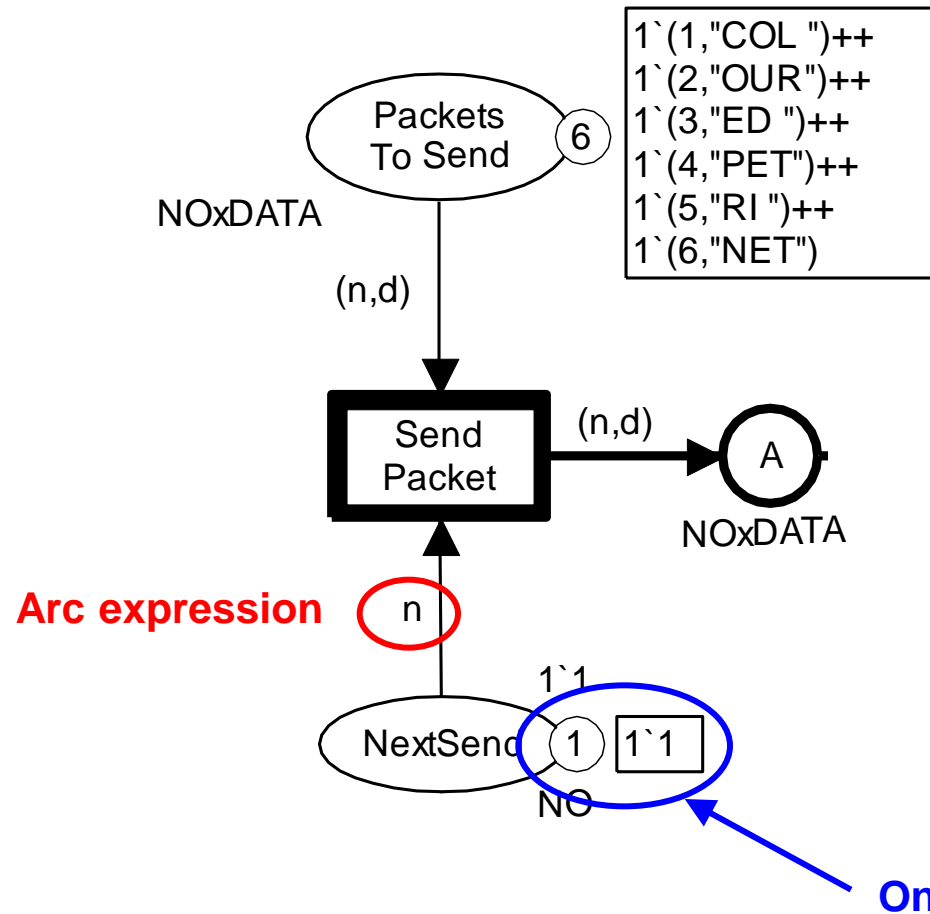
Binding:  $\langle n=? , d=? \rangle$

NO DATA

Transition is **enabled** if we can find a **binding** so that each input arc expressions evaluates to a multi-set of colours that is present on the corresponding input place



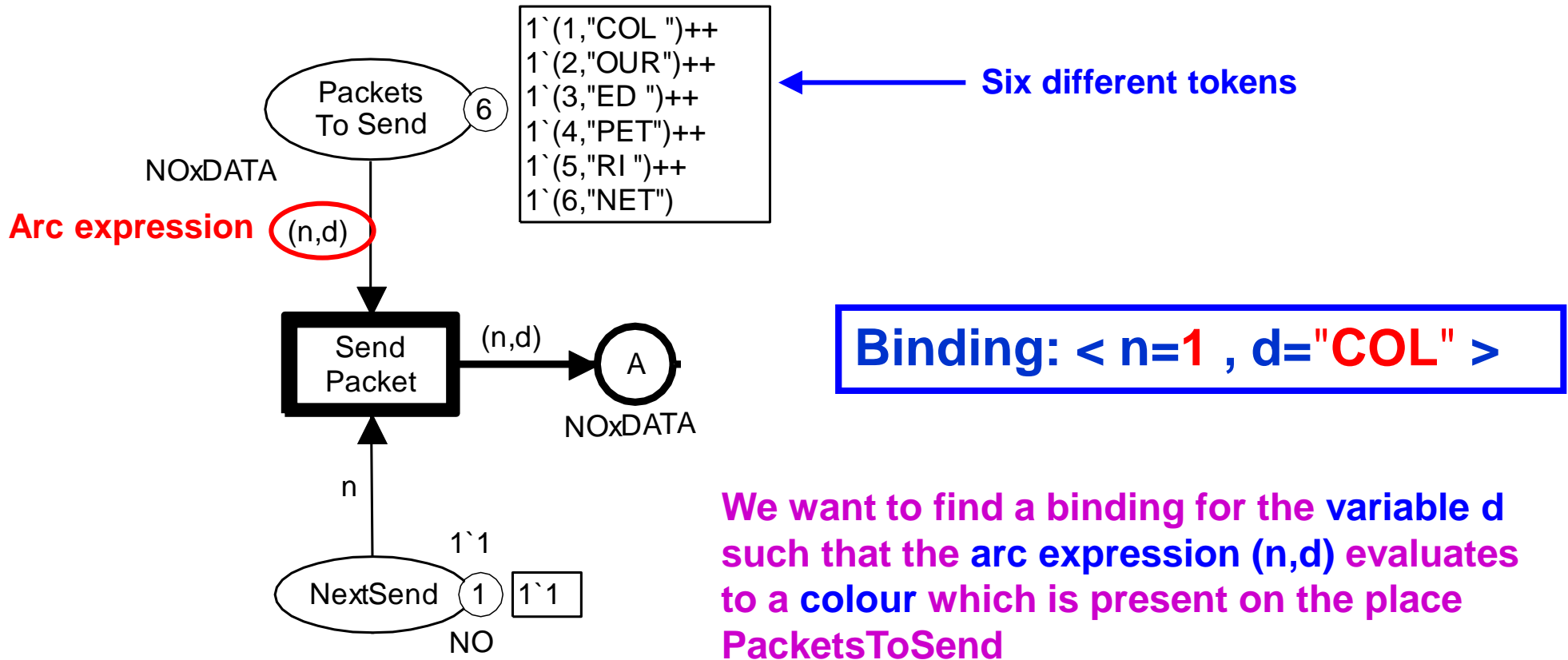
# Enabling of SendPacket



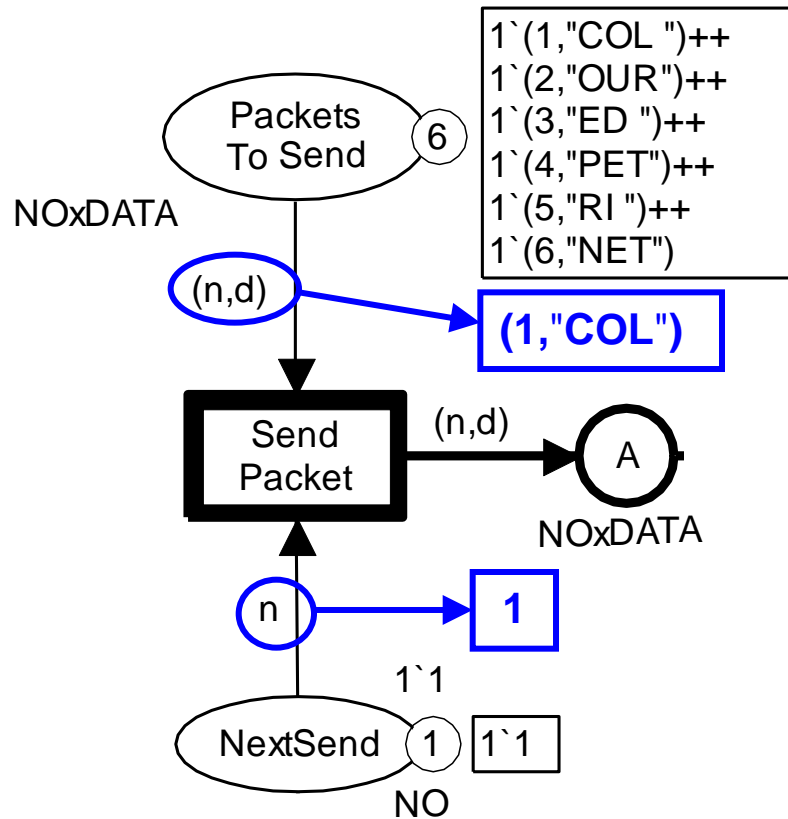
**Binding:  $\langle n=1, d=? \rangle$**

We want to find a binding for the variable  $n$  such that the arc expression  $n$  evaluates to a colour which is present on the place **NextSend**

# Enabling of SendPacket



# Enabling of SendPacket

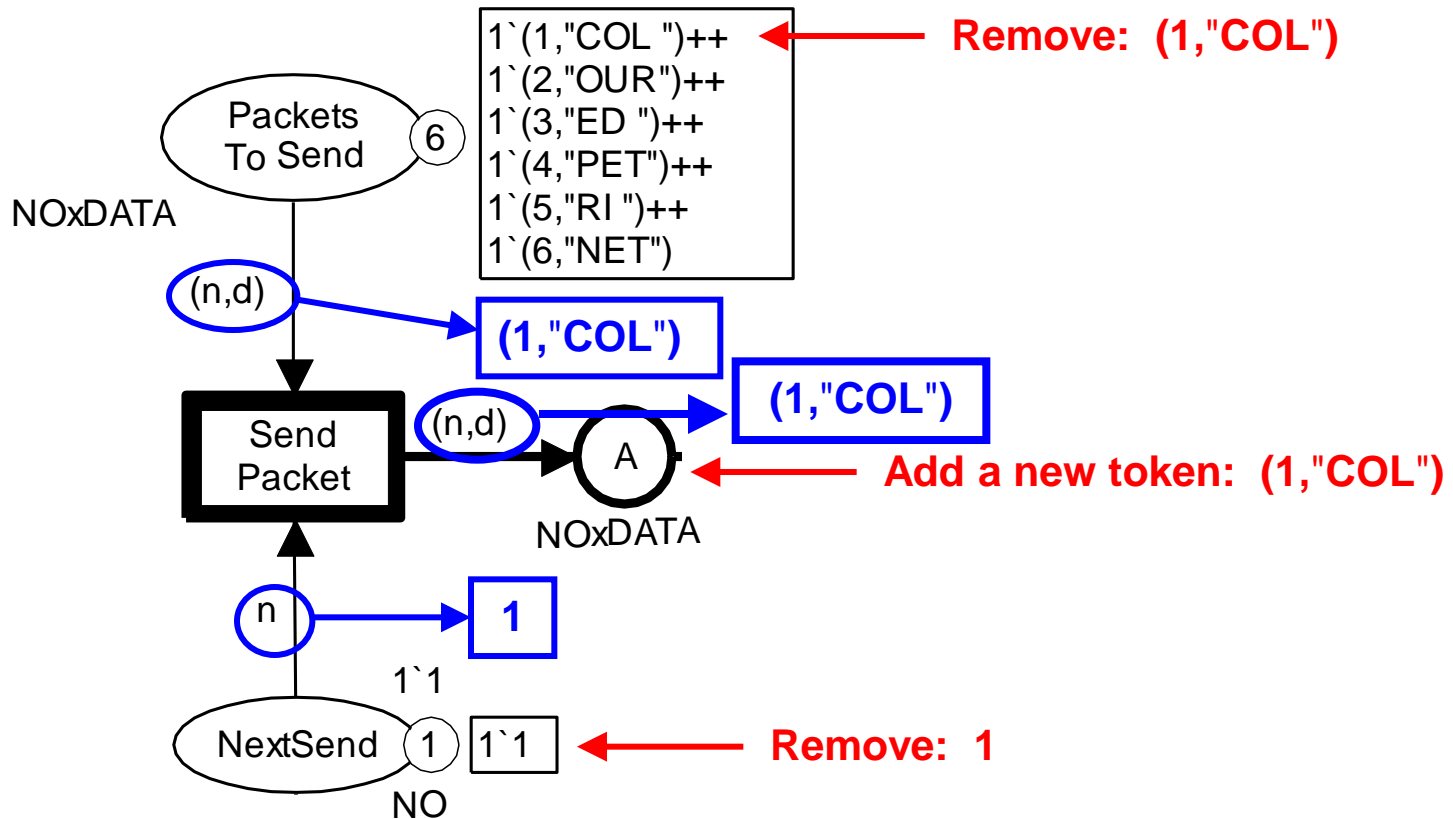


We have found a **binding** so that each **input arc expression** evaluates to a **colour** that is present on the corresponding **input place**

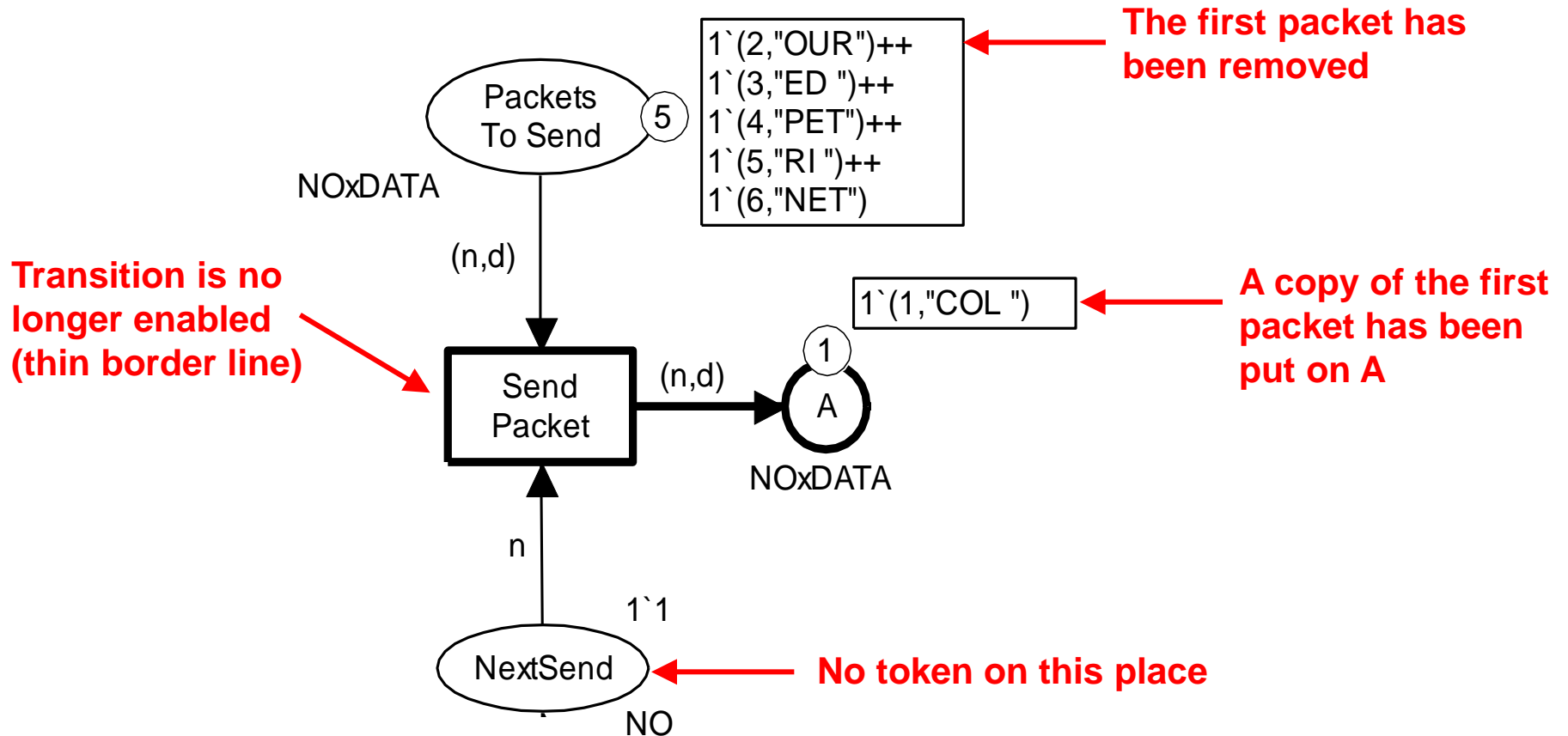
**Binding:  $\langle n=1, d="COL" \rangle$**

**Transition is enabled (ready to occur)**

# Occurrence of SendPacket in binding $\langle n=1, d="COL" \rangle$

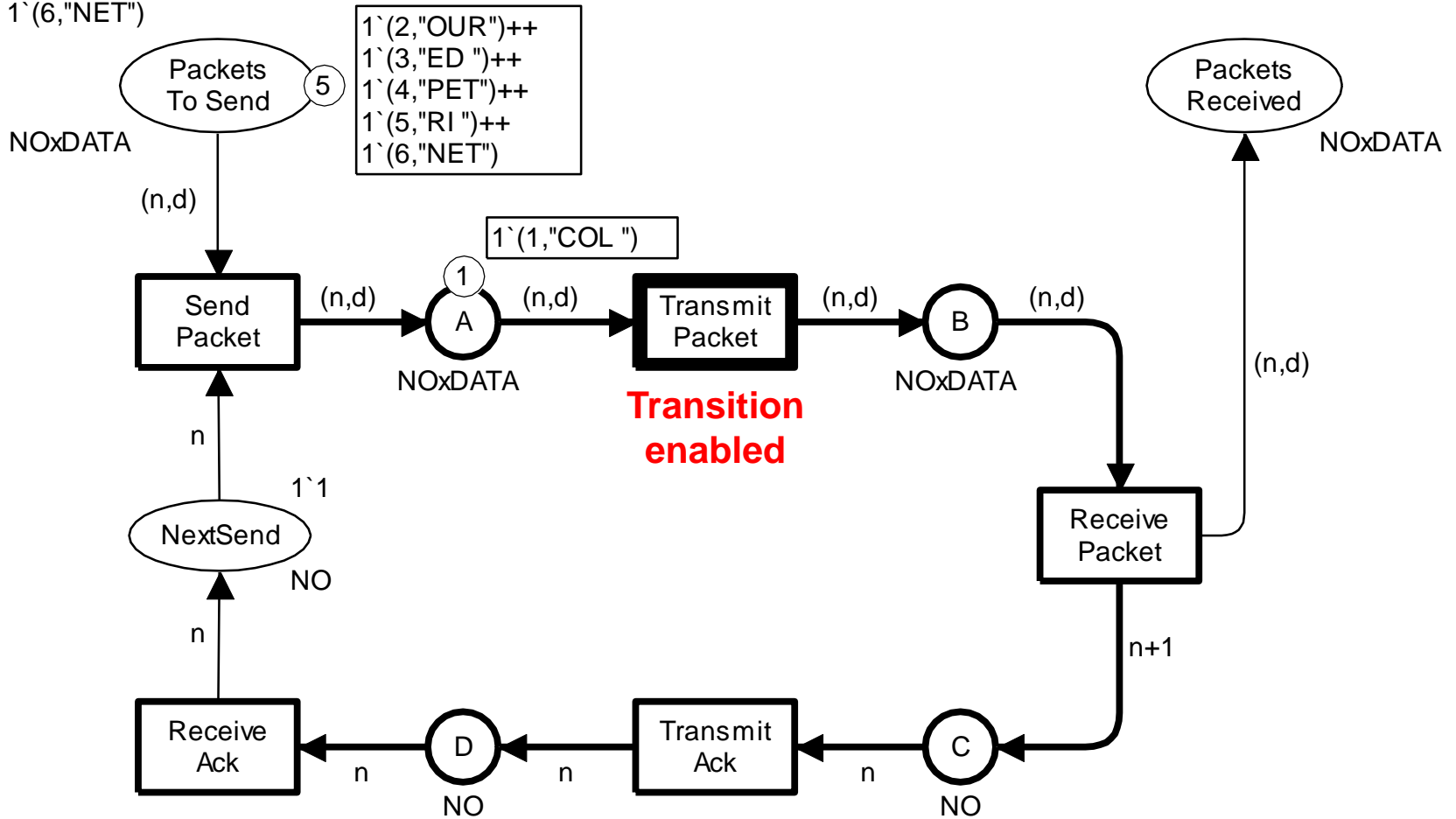


# New marking after occurrence of SendPacket in binding $\langle n=1, d=\text{"COL"} \rangle$

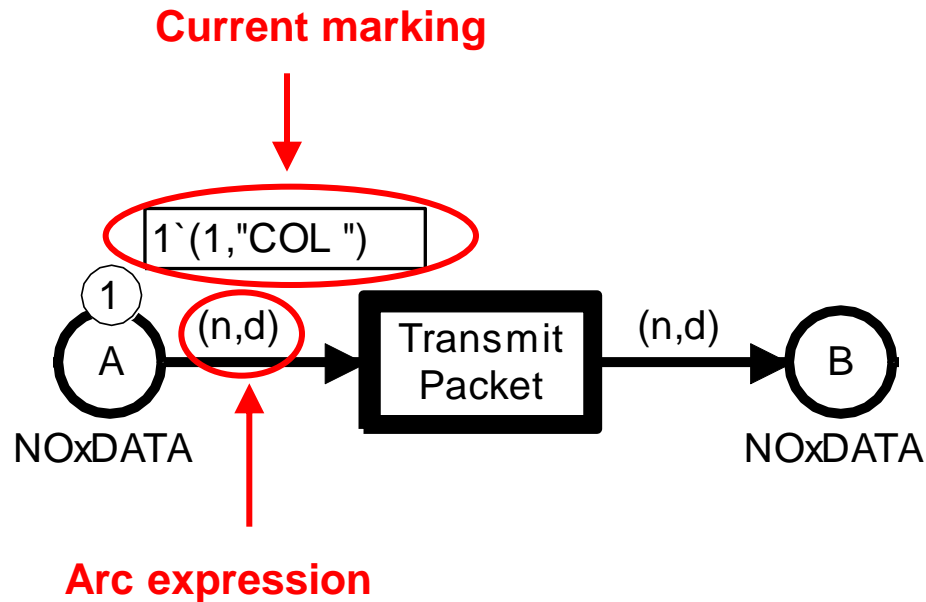


# New marking $M_1$

$1^1(1, "COL ")++$   
 $1^2(2, "OUR")++$   
 $1^3(3, "ED ")++$   
 $1^4(4, "PET")++$   
 $1^5(5, "RI ")++$   
 $1^6(6, "NET")$

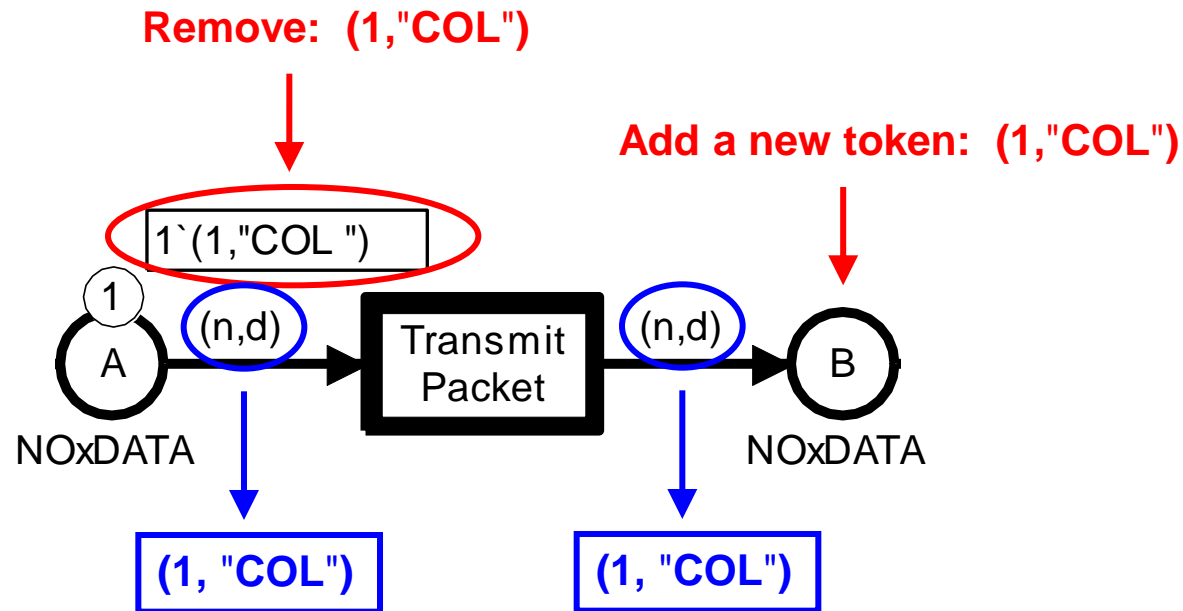


# Binding of TransmitPacket



Binding:  $\langle n=1, d="COL" \rangle$

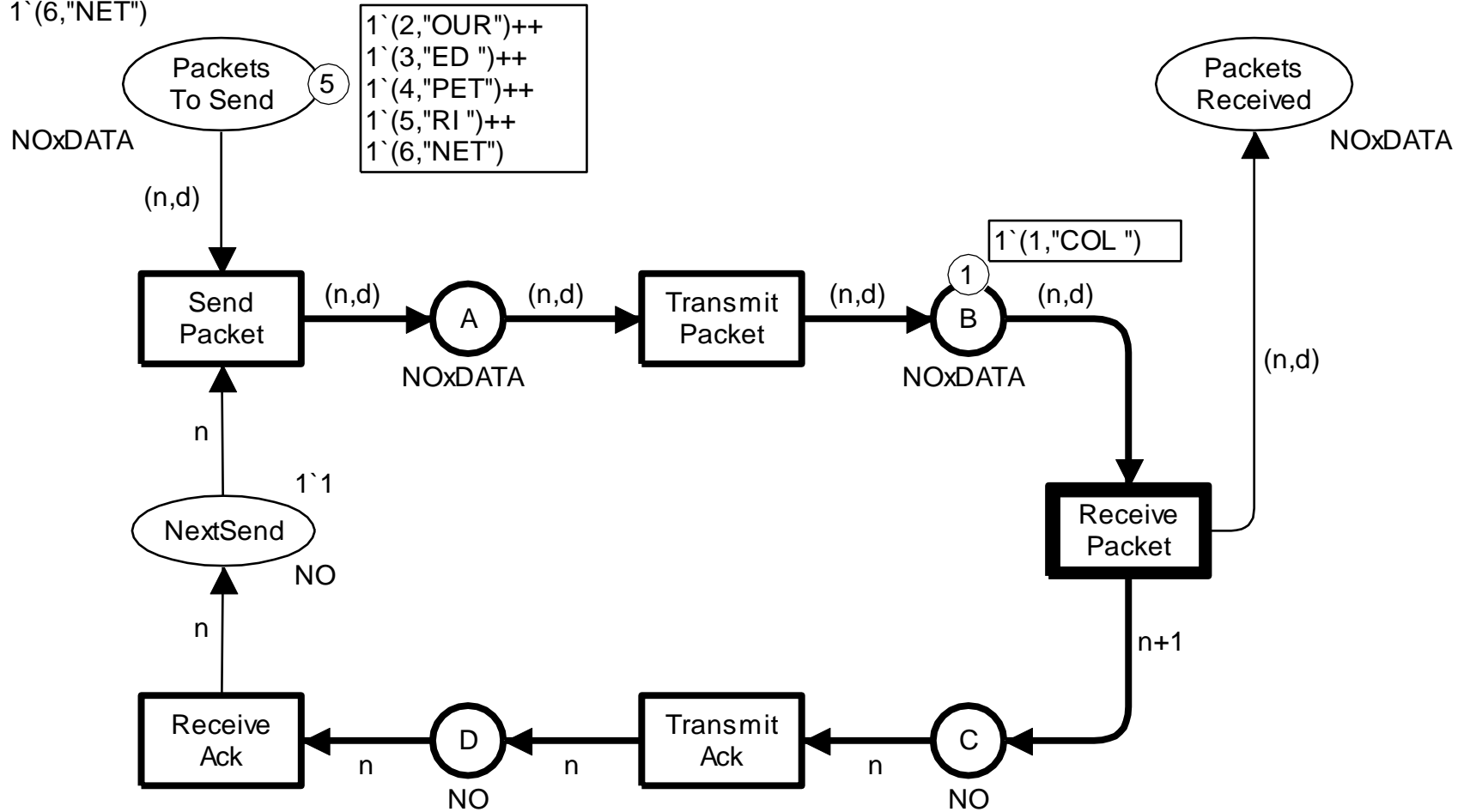
# Occurrence of TransmitPacket in binding $\langle n=1, d="COL" \rangle$





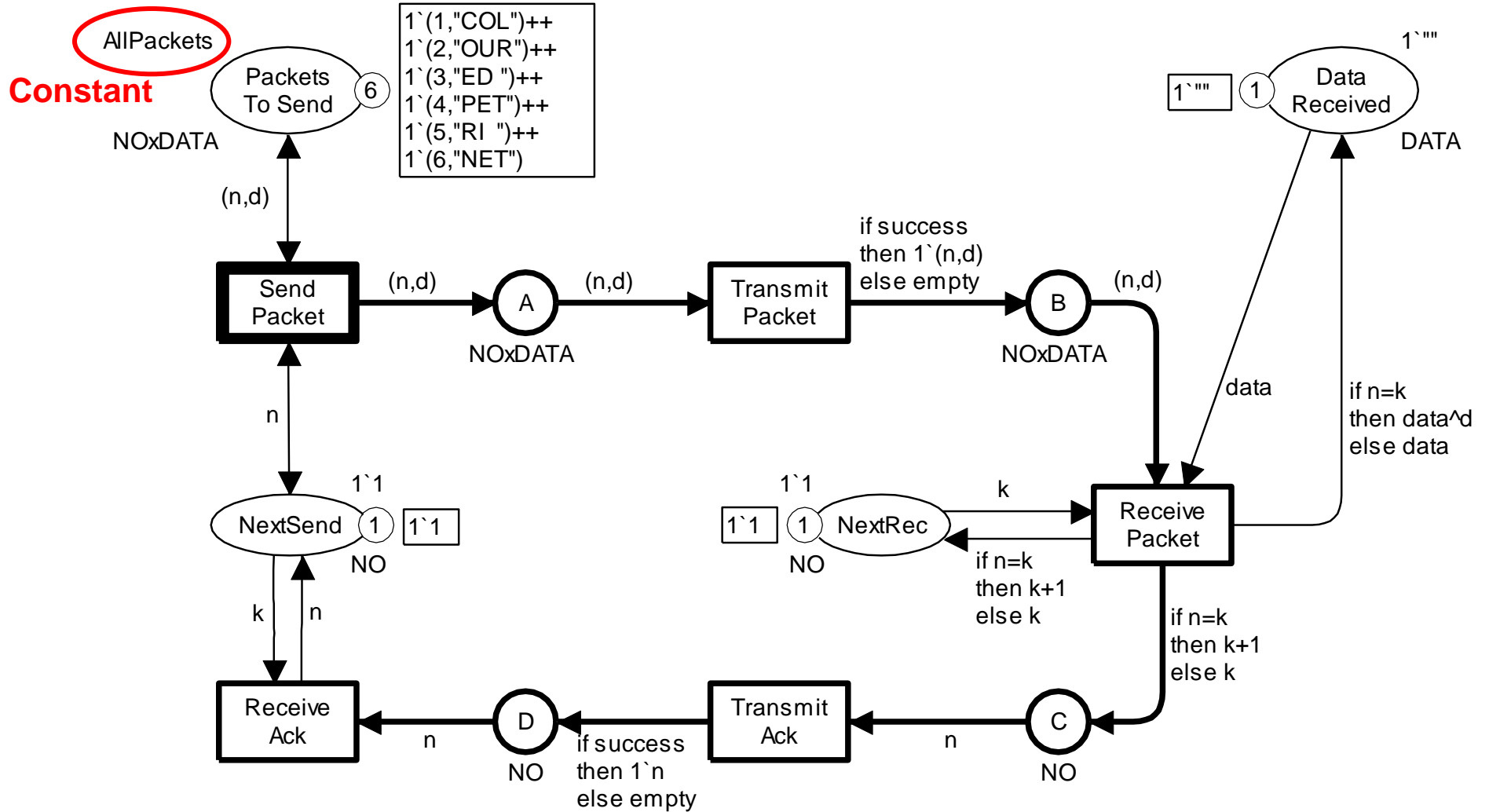
# New marking $M_2$

$1'(1, "COL ")++$   
 $1'(2, "OUR")++$   
 $1'(3, "ED ")++$   
 $1'(4, "PET")++$   
 $1'(5, "RI ")++$   
 $1'(6, "NET")$



# Simulation Demo in CPN Tools

# Second version of protocol



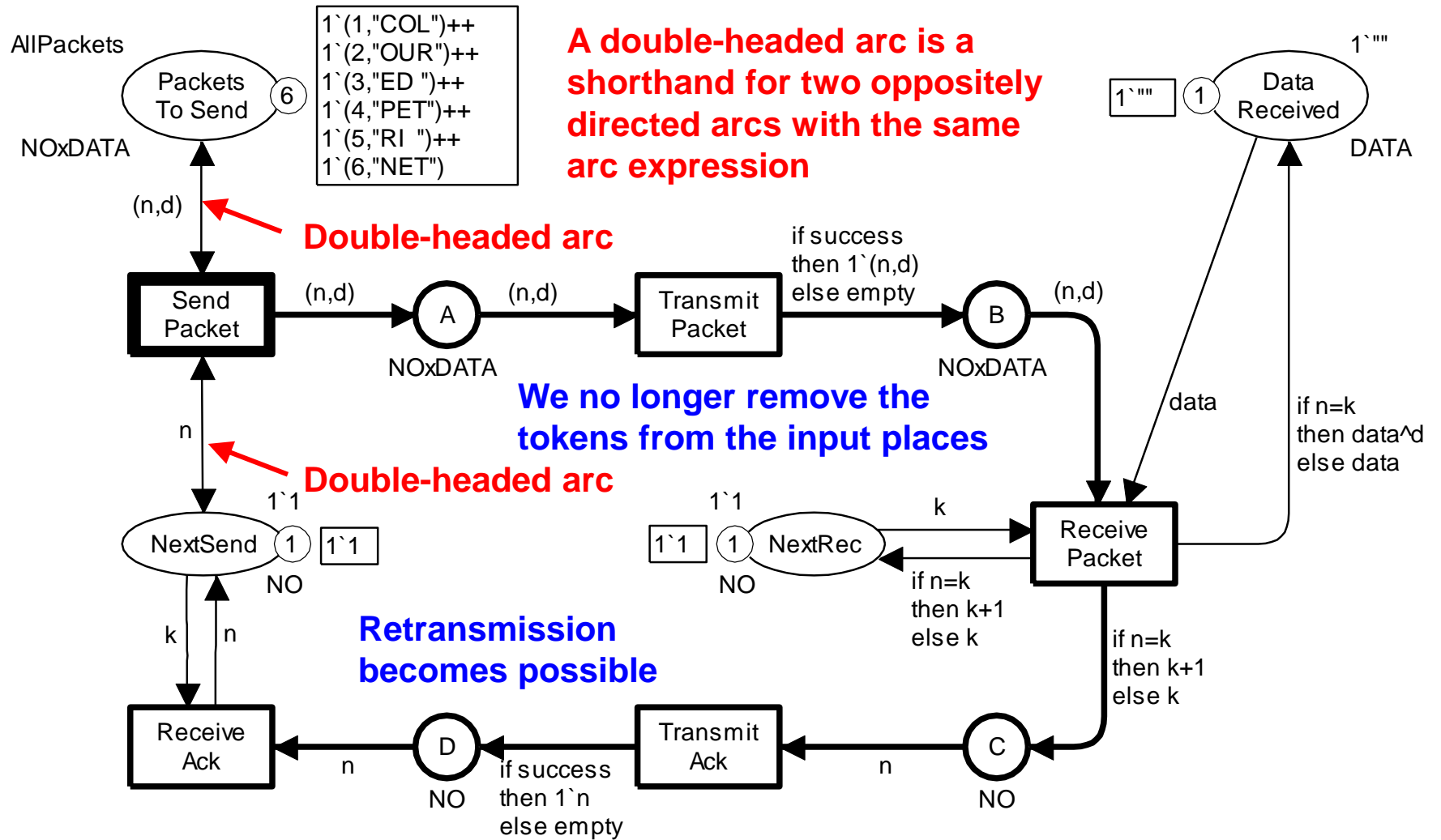
# Declaration of constants

- We use the following **constant** to specify the **initial marking** of PacketsToSend.

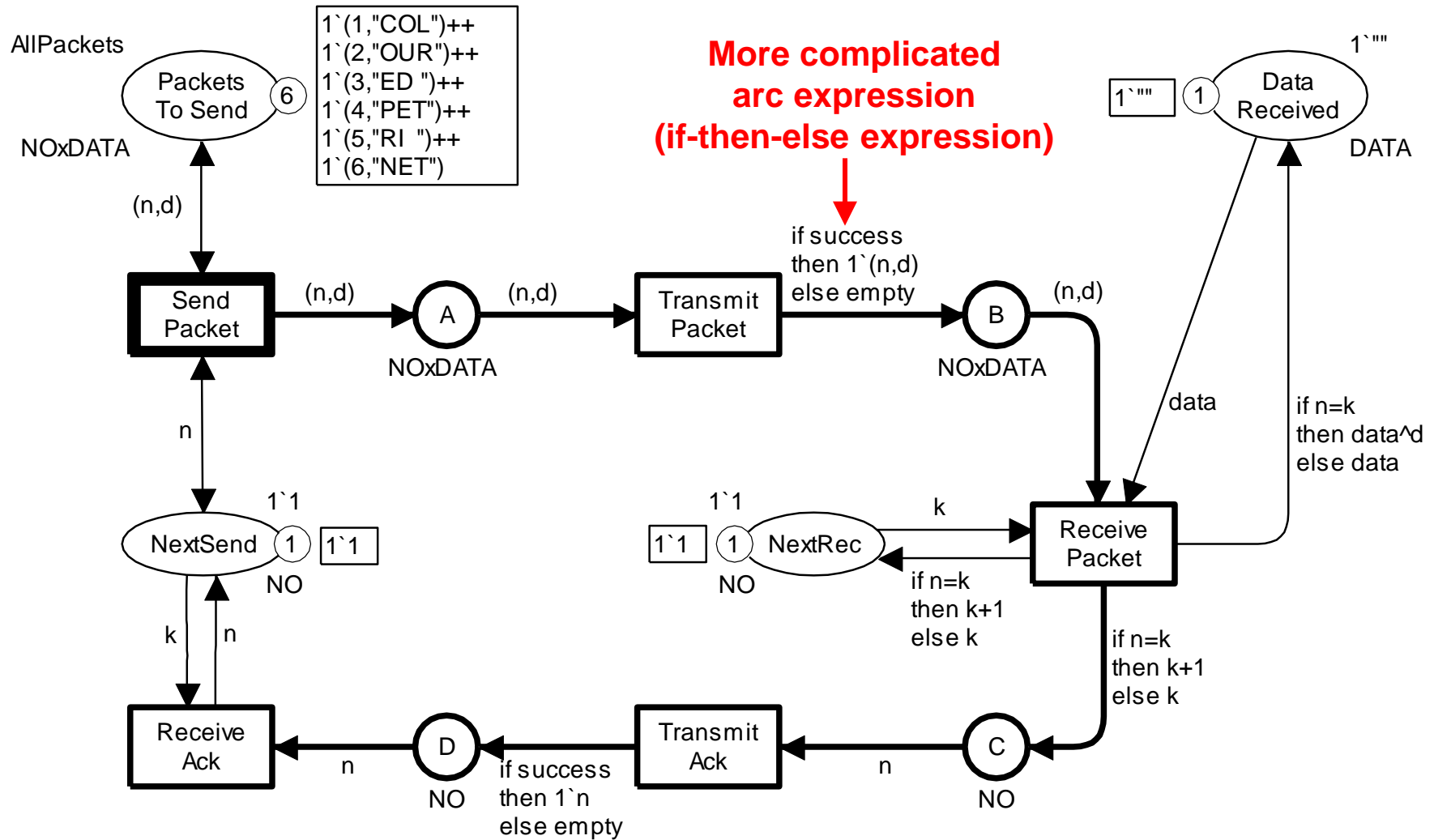
```
val AllPackets = 1 ` (1, "COL") ++ 1 ` (2, "OUR") ++  
                1 ` (3, "ED ") ++ 1 ` (4, "PET") ++  
                1 ` (5, "RI ") ++ 1 ` (6, "NET") ;
```

- Saves a little bit of **space** in the diagram.
- Enhances **readability**.
- Can be **reused** (at other places).

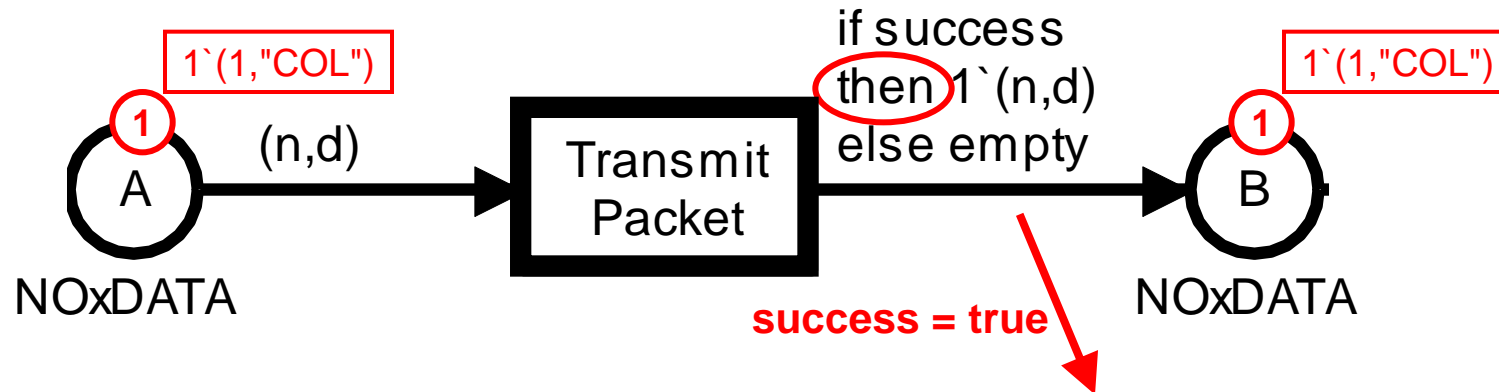
# Double-headed arcs



# More complicated arc expression



# If-then-else expression



New variable:

```
var success : BOOL;
```

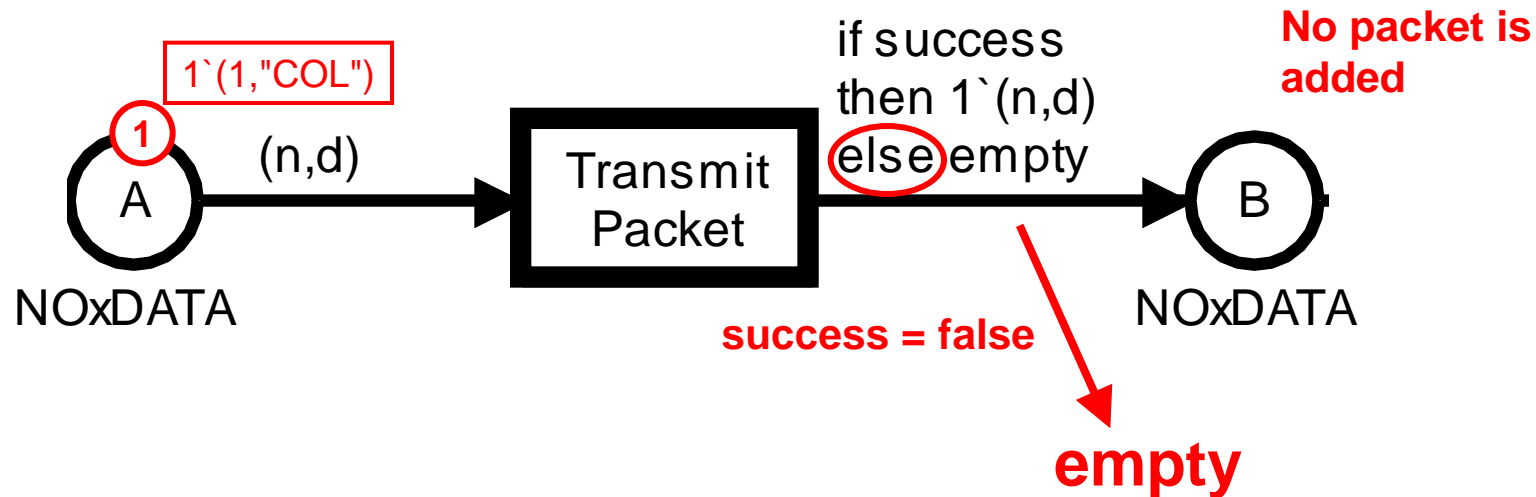
$1^{(1, "COL")}$

Successful transmission  
over the network



```
b+ = <n=1, d="COL", success=true>
b- = <n=1, d="COL", success=false>
```

# If-then-else expression



```
var success : BOOL;
```

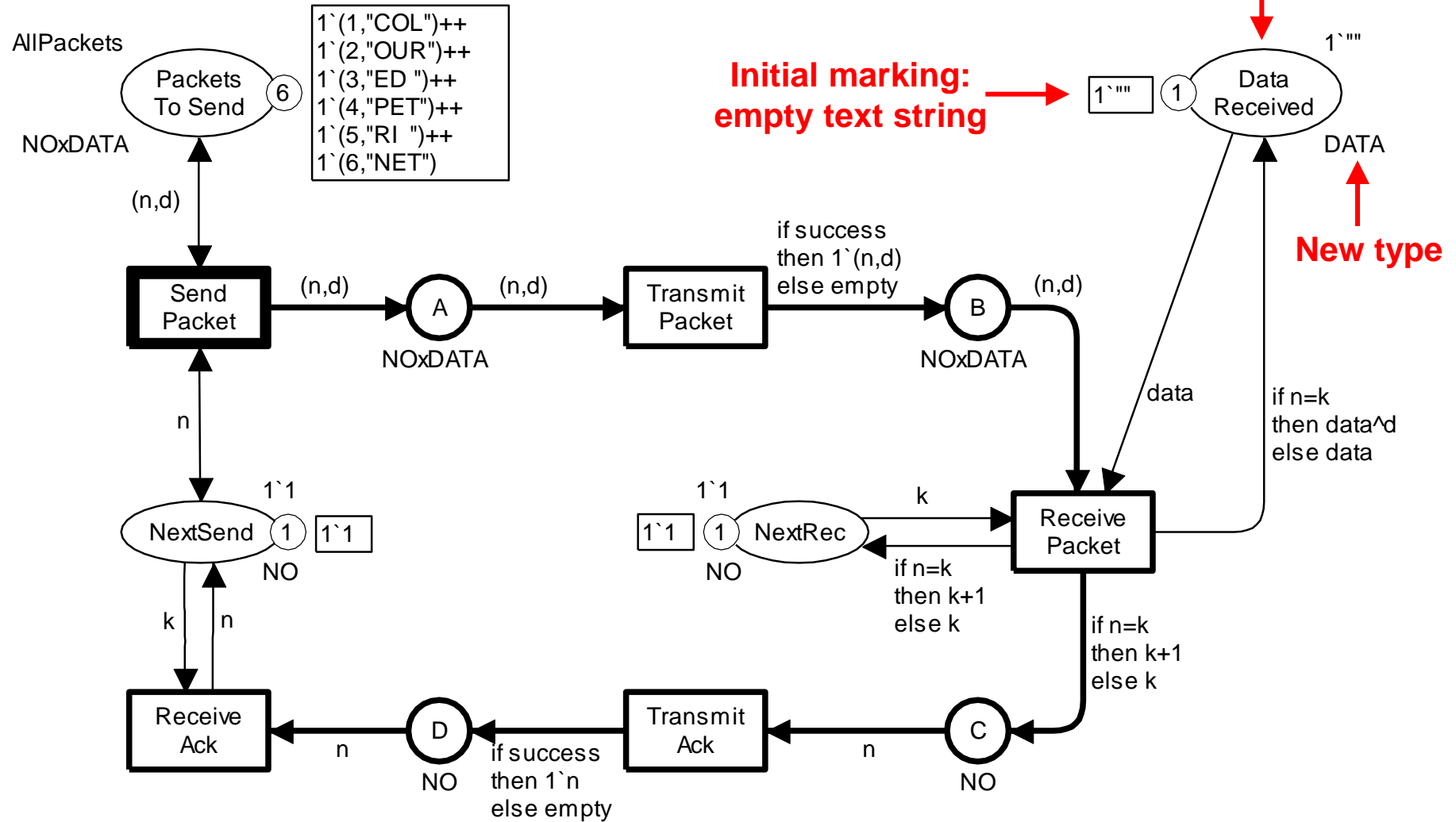
Packet is lost during transmission

→

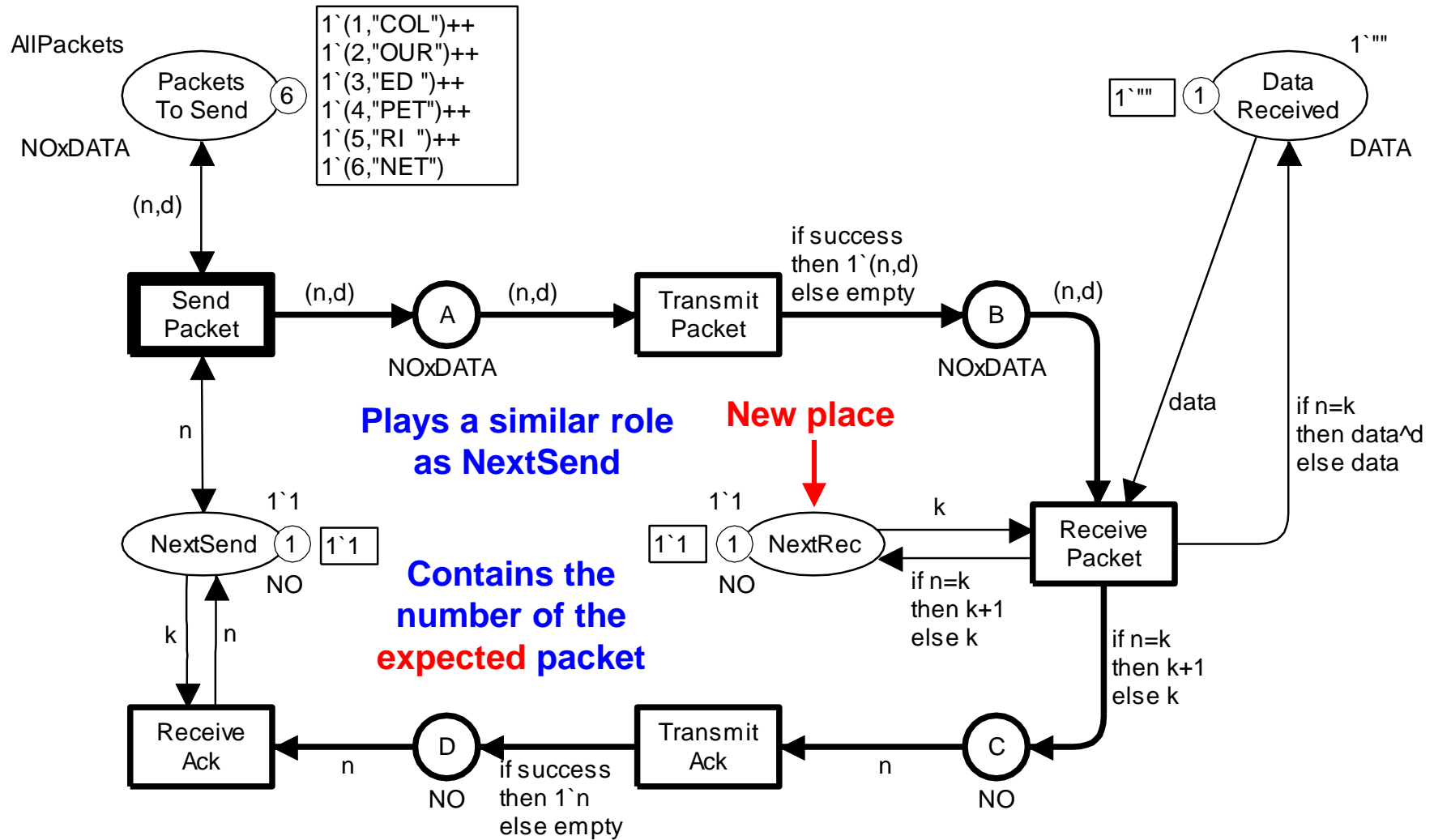
```
b+ = <n=1, d="COL", success=true>
b- = <n=1, d="COL", success=false>
```



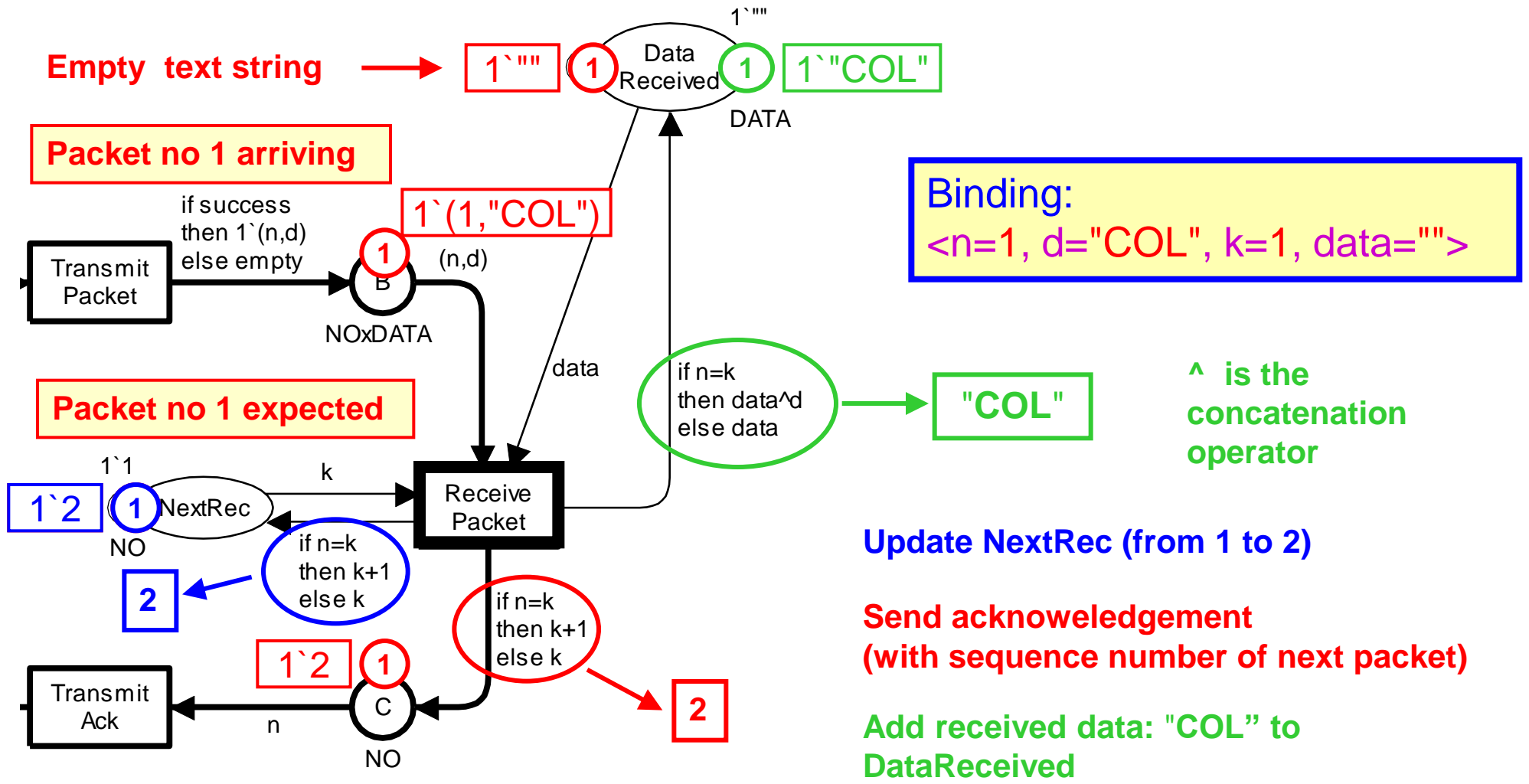
# New name and new type



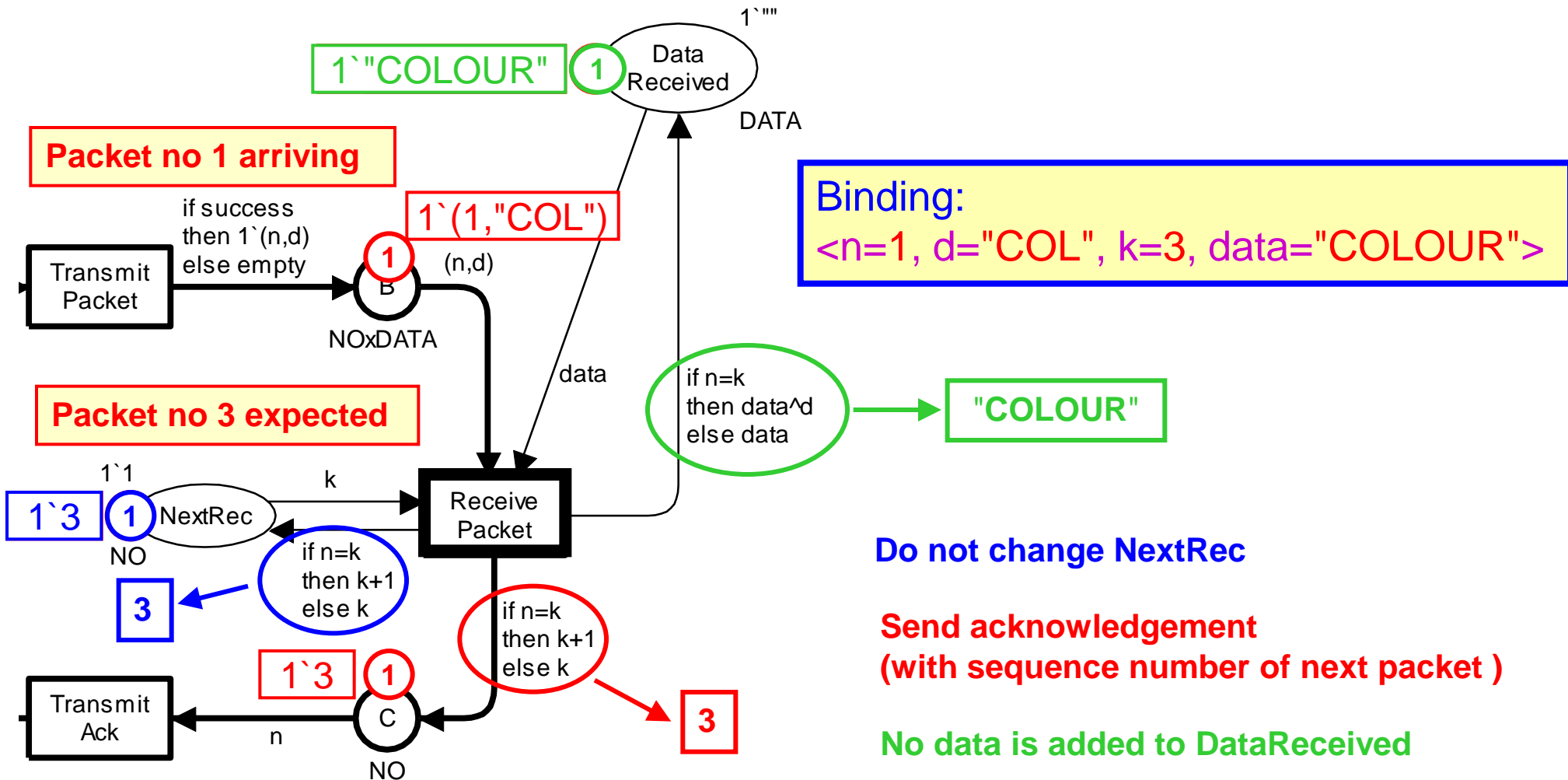
# New place: NextRec



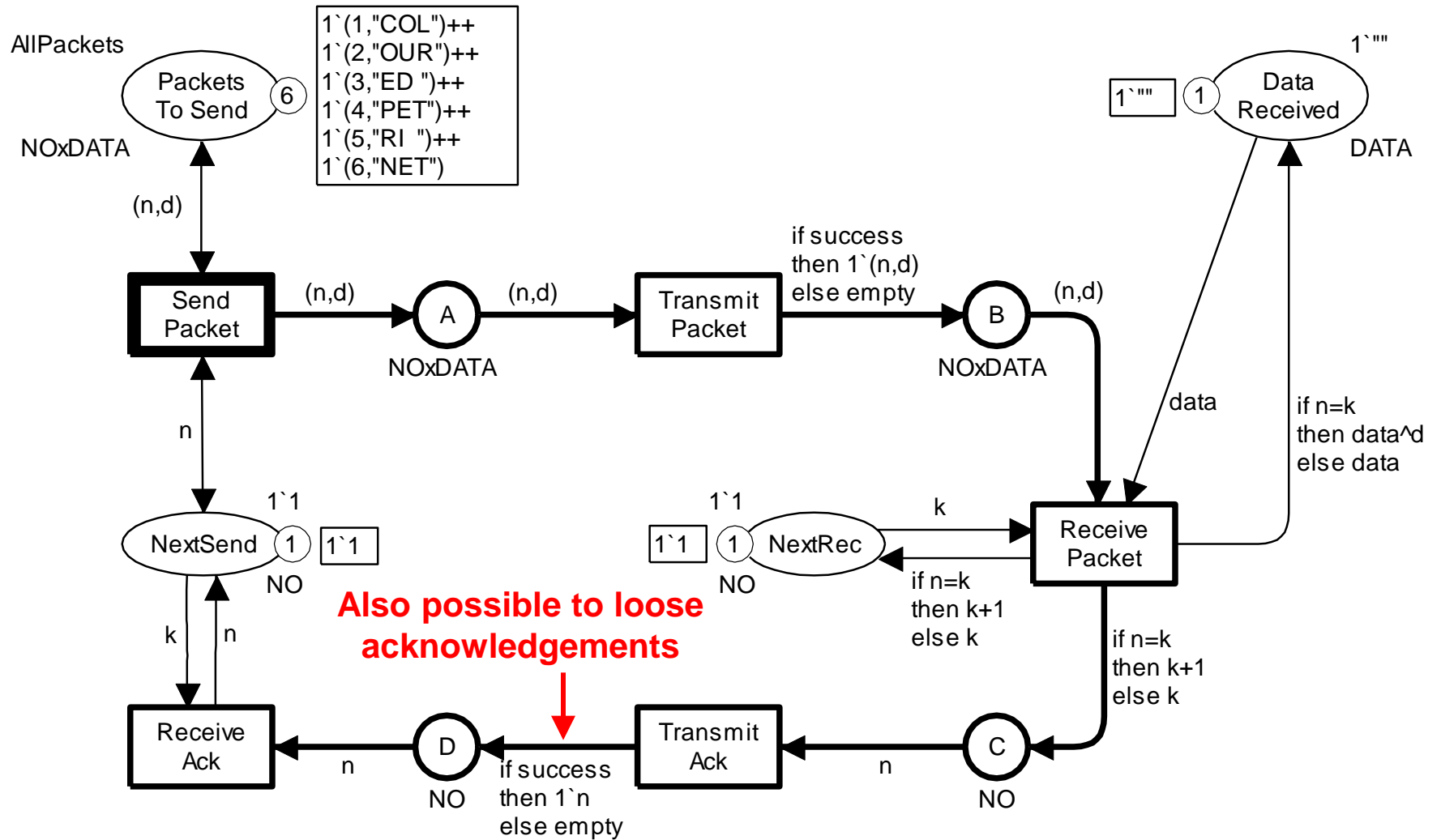
# Correct packet arrives



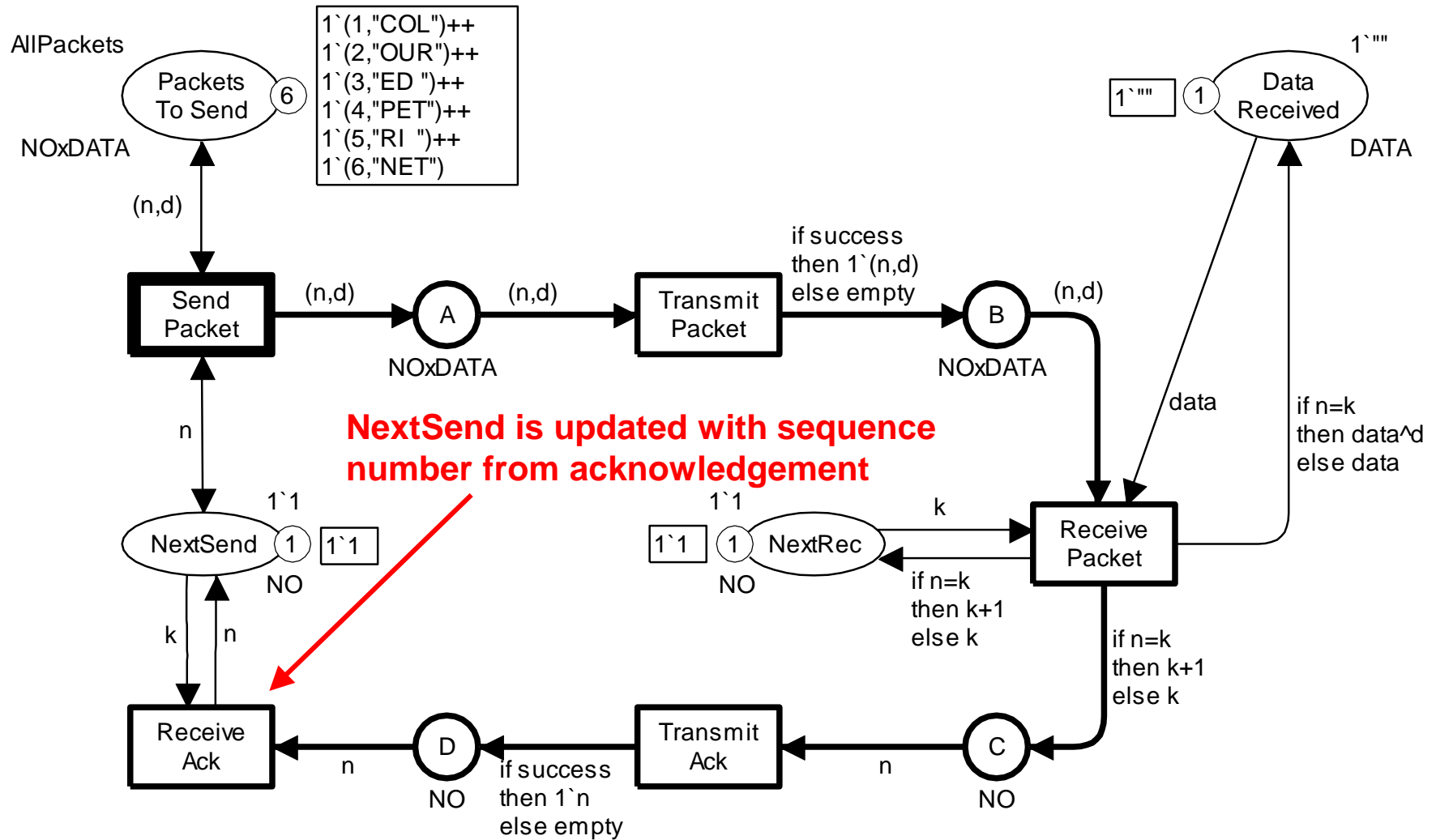
# Wrong packet arrives



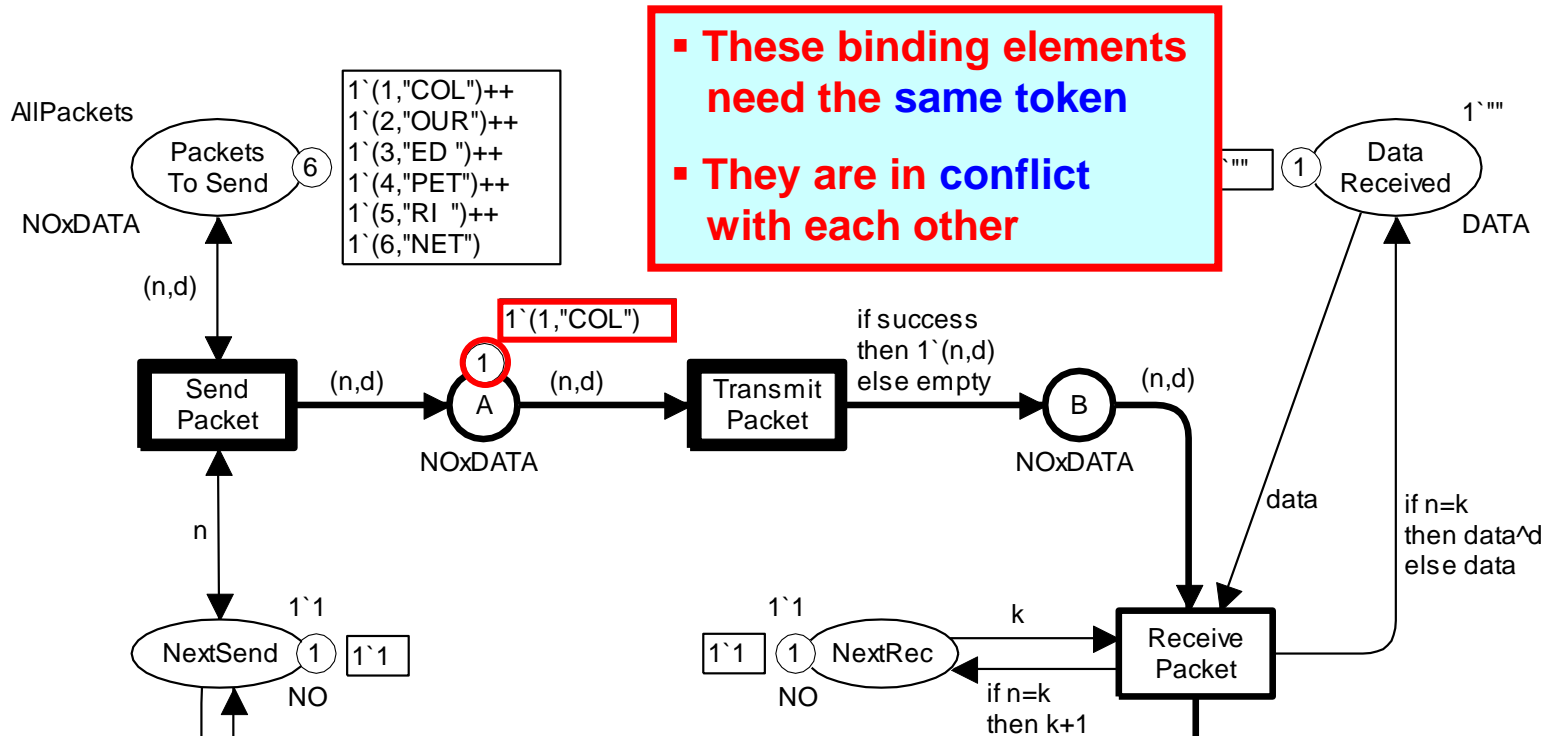
# Acknowledgements can be lost



# NextSend is updated



# Two enabled transitions

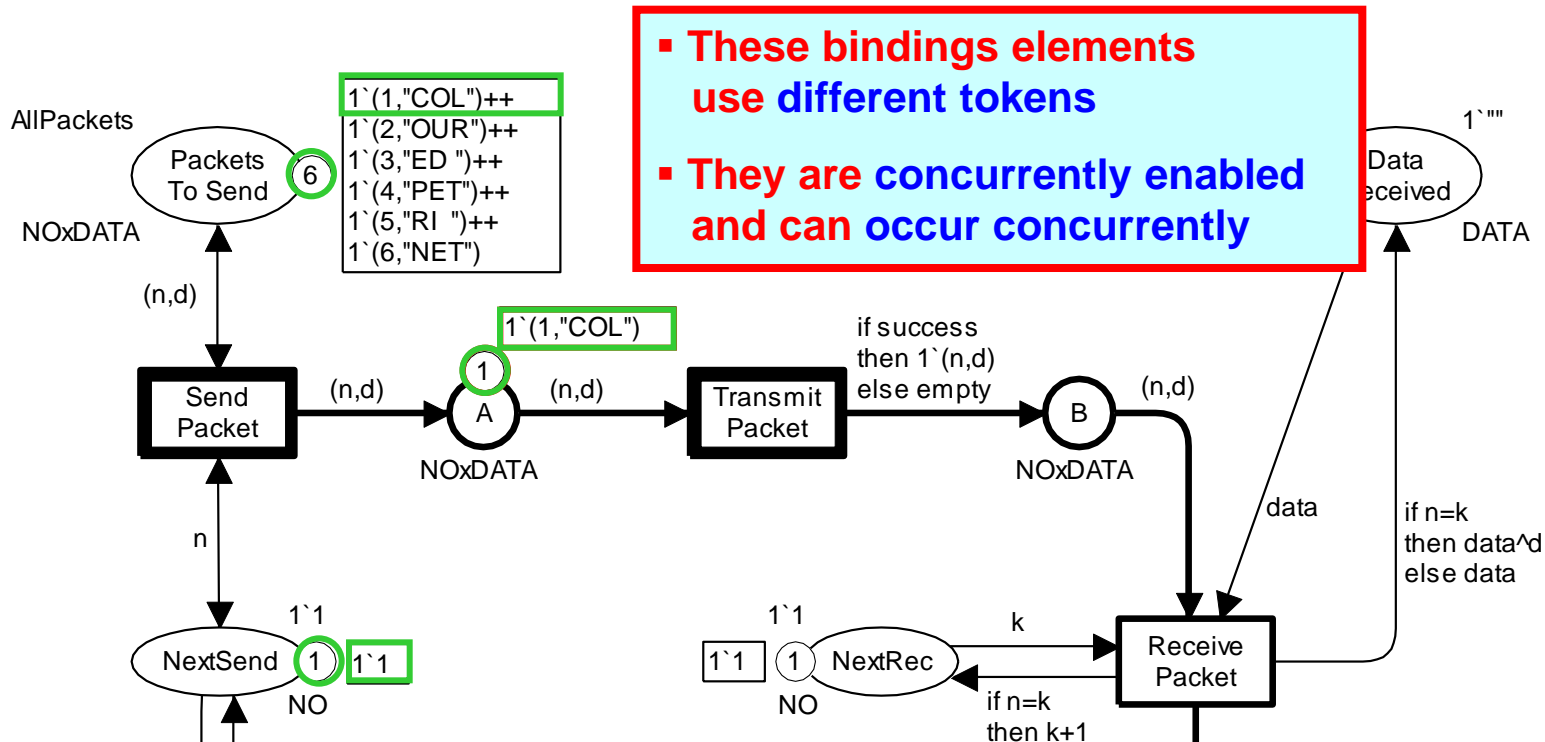


SP = (SendPacket, <n=1, d="COL">)

● TP<sup>+</sup> = (TransmitPacket, <n=1, d="COL", success=true>)

● TP<sup>-</sup> = (TransmitPacket, <n=1, d="COL", success=false>)

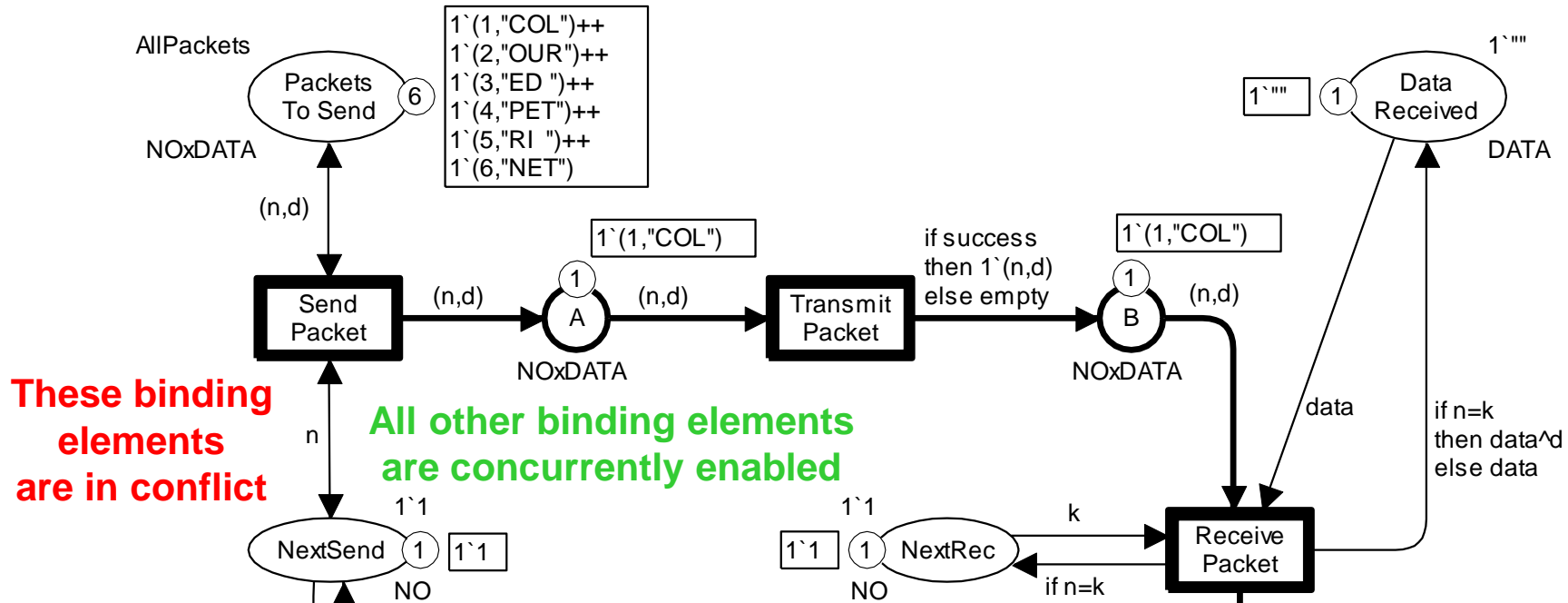
# Two enabled transitions



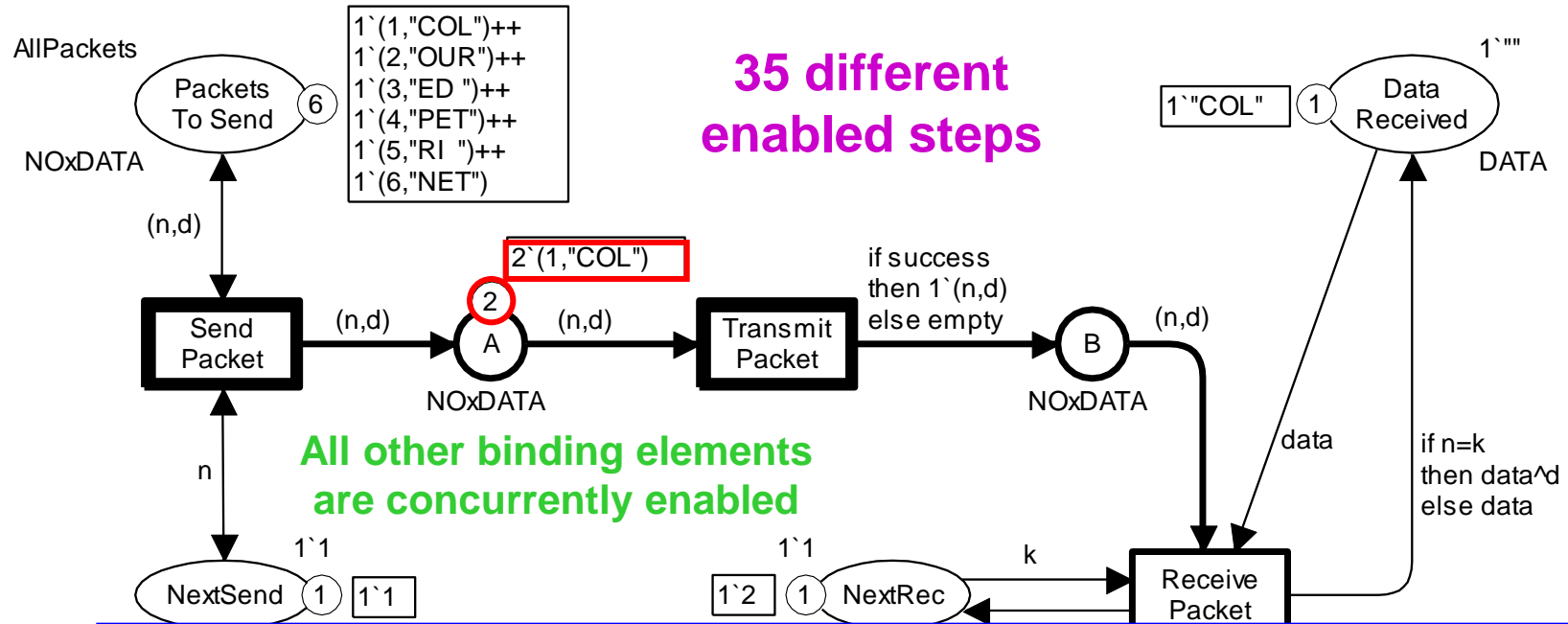
- $SP = (\text{SendPacket}, \langle n=1, d="COL" \rangle)$
- $TP^+ = (\text{TransmitPacket}, \langle n=1, d="COL", \text{success}=\text{true} \rangle)$
- $TP^- = (\text{TransmitPacket}, \langle n=1, d="COL", \text{success}=\text{false} \rangle)$



# Three enabled transitions



# Three enabled transitions

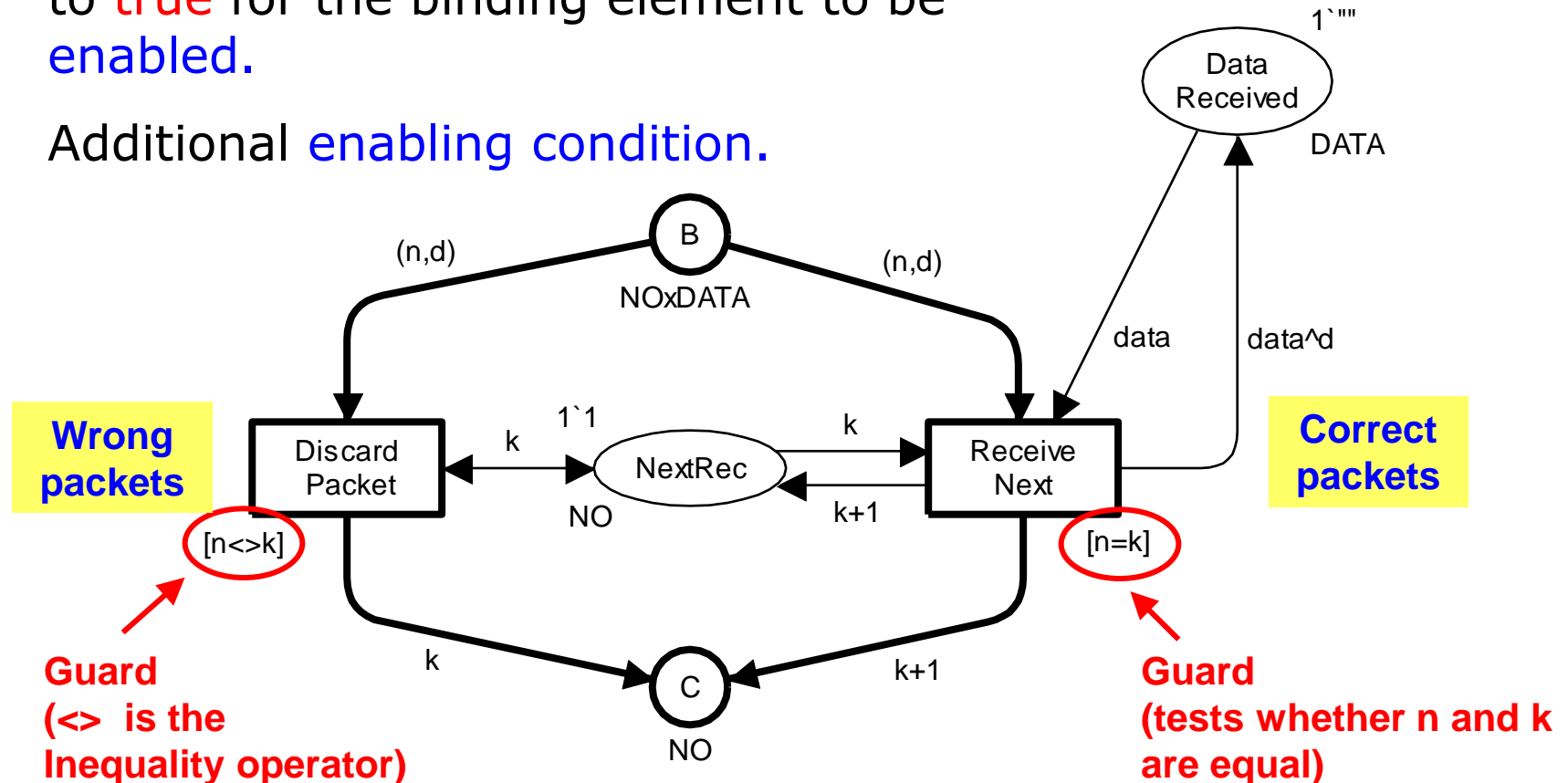


- SP = (SendPacket,  $\langle n=1, d=\text{"COL"} \rangle$ )
- TP<sup>+</sup> = (TransmitPacket,  $\langle n=1, d=\text{"COL"}, \text{success}=\text{true} \rangle$ )
- TP<sup>-</sup> = (TransmitPacket,  $\langle n=1, d=\text{"COL"}, \text{success}=\text{false} \rangle$ )
- TA<sup>+</sup> = (TransmitAck,  $\langle n=2, \text{success}=\text{true} \rangle$ )
- TA<sup>-</sup> = (TransmitAck,  $\langle n=2, \text{success}=\text{false} \rangle$ )

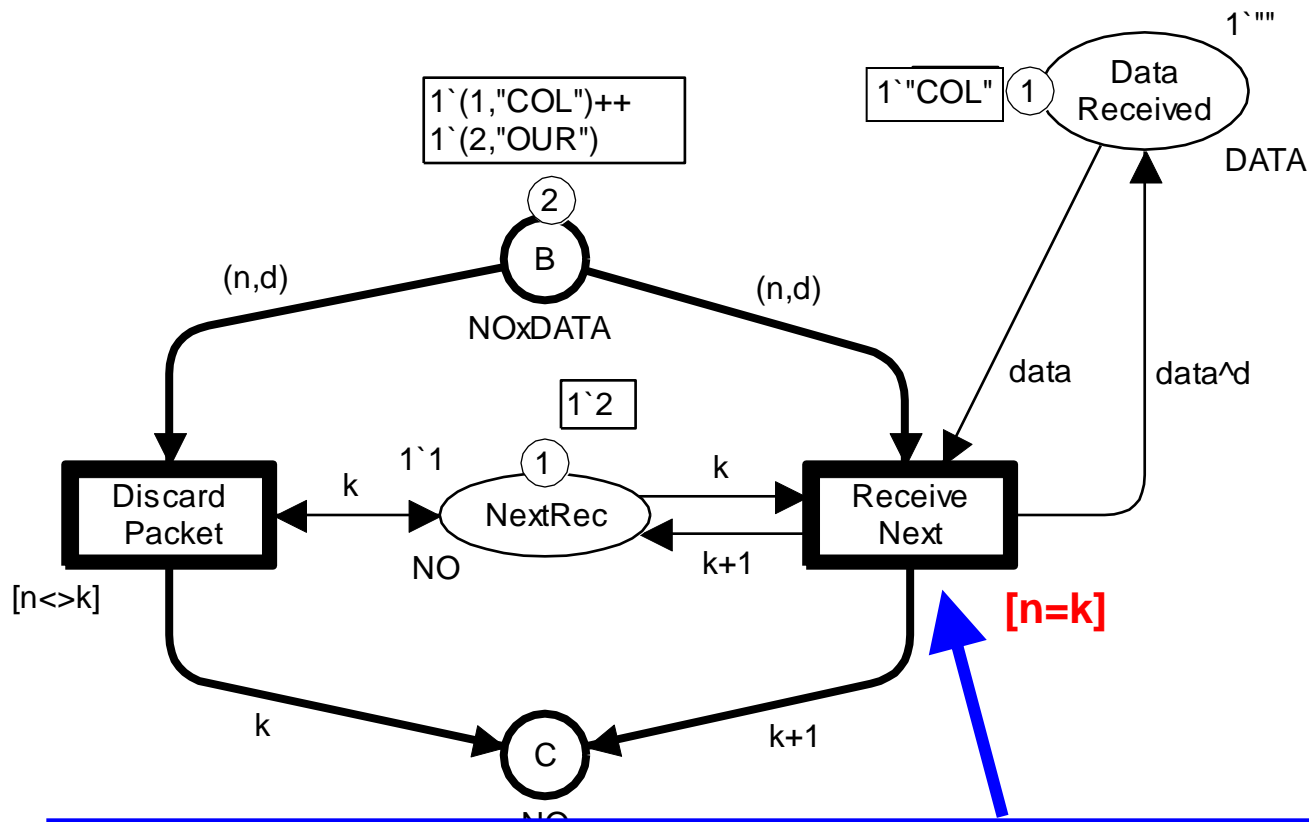
# Simulation Demo in CPN Tools

# Transitions can have a guard

- Boolean expression which must evaluate to **true** for the binding element to be enabled.
- Additional **enabling condition**.



# Guard must evaluate to true



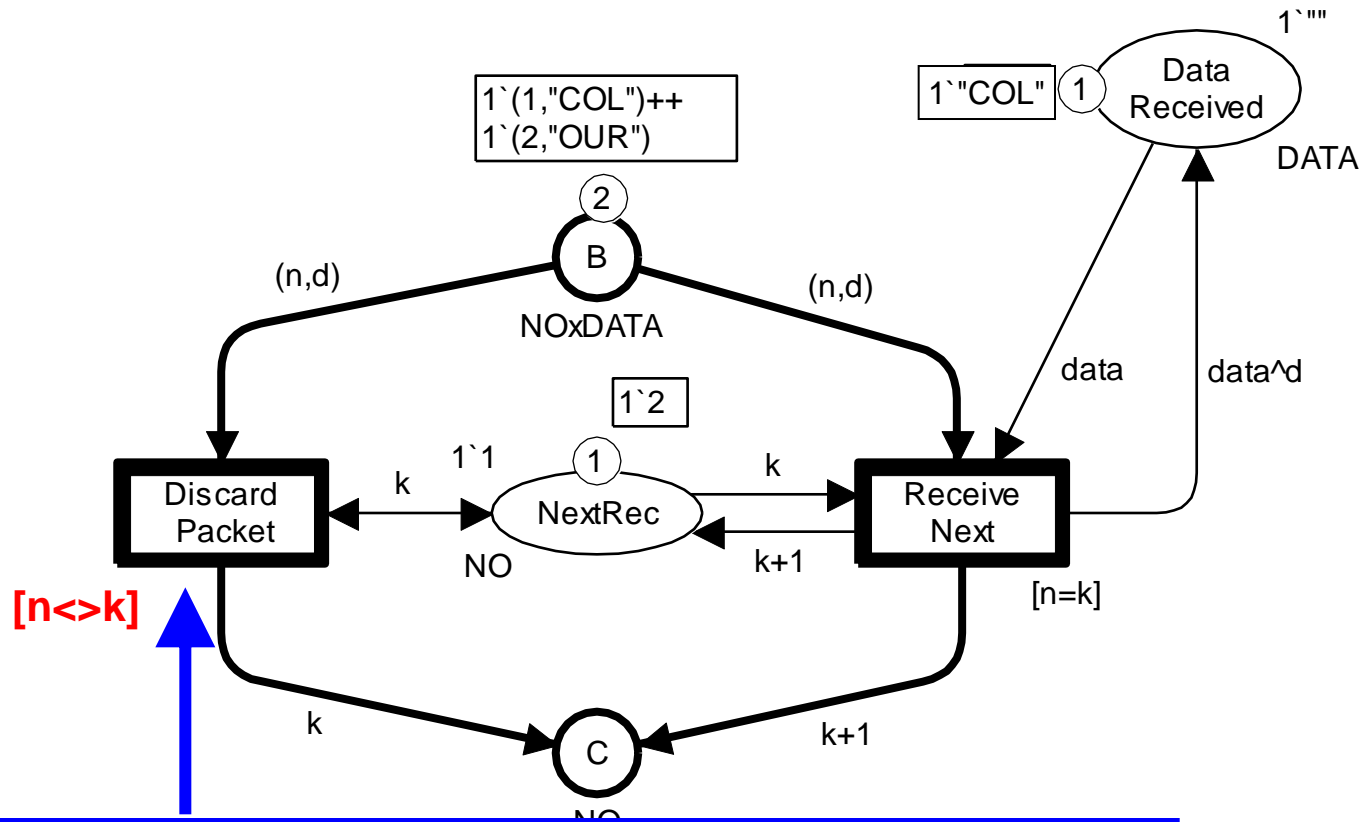
false

$RN_1 = (\text{ReceiveNext}, \langle n=1, k=2, d=\text{"COL"}, \text{data}=\text{"COL"} \rangle)$

true

$RN_2 = (\text{ReceiveNext}, \langle n=2, k=2, d=\text{"OUR"}, \text{data}=\text{"COL"} \rangle)$

# Guard must evaluate to true



true  
false

$DP_1 = (\text{DiscardPacket}, \langle n=1, k=2, d=\text{"COL"} \rangle)$   
 $DP_2 = (\text{DiscardPacket}, \langle n=2, k=2, d=\text{"OUR"} \rangle)$

# Editing Demo in CPN Tools

# Questions





# Assignment 3 – Task 1

- In the lectures you have seen a stop-and-wait protocol. The sender keeps sending the same packet until a matching acknowledgement is received. In a sliding window protocol it is possible for the sender to transmit several packets to the receiver before receiving an acknowledgement. The sender has a window containing a number of data packets which are currently under transmission and for which acknowledgements have not yet been received.

## **Task 1: Modify the CPN model such that it models a Go-Back-N Window Protocol.**

- In a Go-Back-N protocol, the sender sends all data packets in the current window and then waits for acknowledgements. If no acknowledgement is received (within a certain amount of time), the data packets in the window are all retransmitted.
- The CPN model of the stop-and-wait protocol can be downloaded from:  
<http://www.cs.au.dk/~cpnbook/models/chapter2/2-10NondeterministicProtocol.cpn>
- It should only be necessary to modify the Sender part of the CPN model. You should change the colour set and the initial marking of the NextSend place, so that it contains information about the start and end of the window. Having done this you should change the arc inscriptions of the surrounding arcs so that they implement a sliding windows strategy.
- Use simulation to validate the correctness of your protocol. It may be useful to change the initial marking of PacketToSend so that you get more packets to work with.